THE DATASHEET OF

AT90S1200-4PC

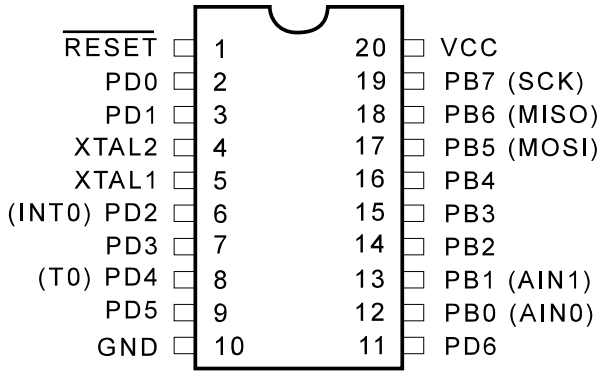
www.win-source.net 0086-755-83957316

Features

- Utilizes the AVR® RISC Architecture
- AVR – High-performance and Low-power RISC Architecture
  - 89 Powerful Instructions – Most Single Clock Cycle Execution
  - 32 x 8 General Purpose Working Registers
  - Up to 12 MIPS Throughput at 12 MHz
- Data and Non-volatile Program Memory
  - 1K Byte of In-System Programmable Flash
    Endurance: 1,000 Write/Erase Cycles
  - 64 Bytes of In-System Programmable EEPROM
    Endurance: 100,000 Write/Erase Cycles
  - Programming Lock for Flash Program and EEPROM Data Security
- Peripheral Features
  - One 8-bit Timer/Counter with Separate Prescaler
  - On-chip Analog Comparator
  - Programmable Watchdog Timer with On-chip Oscillator
  - SPI Serial Interface for In-System Programming
- Special Microcontroller Features
  - Low-power Idle and Power-down Modes
  - External and Internal Interrupt Sources
  - Selectable On-chip RC Oscillator for Zero External Components
- Specifications
  - Low-power, High-speed CMOS Process Technology
  - Fully Static Operation
- Power Consumption at 4 MHz, 3V, 25°C
  - Active: 2.0 mA
  - Idle Mode: 0.4 mA
  - Power-down Mode: <1 µA
- I/O and Packages
  - 15 Programmable I/O Lines
  - 20-pin PDIP, SOIC and SSOP
- Operating Voltages
  - 2.7 - 6.0V (AT90S1200-4)
  - 4.0 - 6.0V (AT90S1200-12)
- Speed Grades
  - 0 - 4 MHz, (AT90S1200-4)
  - 0 - 12 MHz, (AT90S1200-12)

Pin Configuration

PDIP/SOIC/SSOP



8-bit AVR® Microcontroller with 1K Byte of In-System Programmable Flash

AT90S1200

Rev. 0838H–AVR–03/02

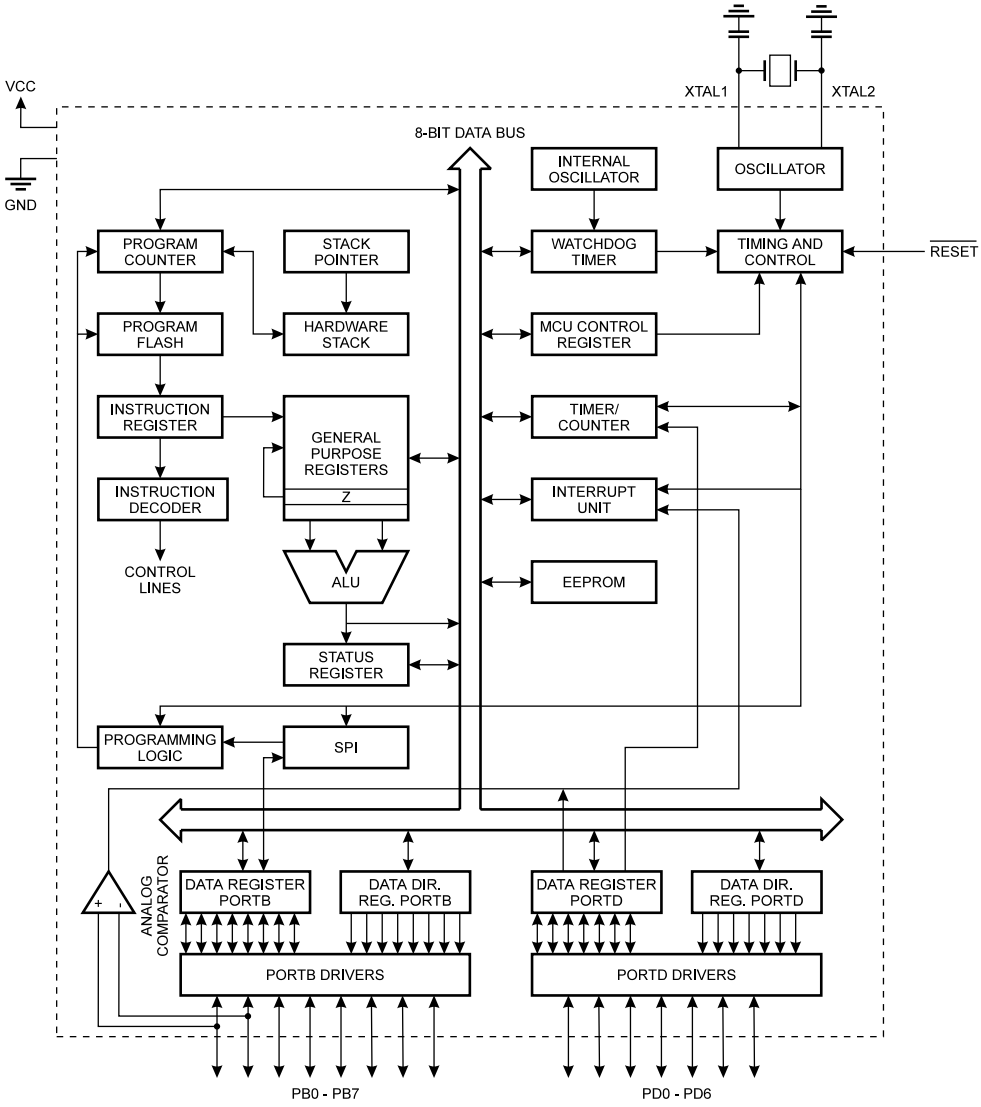
Description

The AT90S1200 is a low-power CMOS 8-bit microcontroller based on the AVR RISC architecture. By executing powerful instructions in a single clock cycle, the AT90S1200 achieves throughputs approaching 1 MIPS per MHz allowing the system designer to optimize power consumption versus processing speed.

The AVR core combines a rich instruction set with the 32 general purpose working registers. All the 32 registers are directly connected to the Arithmetic Logic Unit (ALU), allowing two independent registers to be accessed in one single instruction executed in one clock cycle. The resulting architecture is more code efficient while achieving throughputs up to ten times faster than conventional CISC microcontrollers.

Block Diagram

Figure 1. The AT90S1200 Block Diagram



The architecture supports high-level languages efficiently as well as extremely dense assembler code programs. The AT90S1200 provides the following features: 1K byte of In-System Programmable Flash, 64 bytes EEPROM, 15 general purpose I/O lines, 32 general purpose working registers, internal and external interrupts, programmable watchdog timer with internal oscillator, an SPI serial port for program downloading and two software selectable power-saving modes. The Idle Mode stops the CPU while allow-

ing the Registers, Timer/Counter, Watchdog and Interrupt system to continue functioning. The Power-down mode saves the register contents but freezes the Oscillator, disabling all other chip functions until the next External Interrupt or hardware Reset.

The device is manufactured using Atmel's high-density nonvolatile memory technology. The On-chip In-System Programmable Flash allows the program memory to be reprogrammed in-system through an SPI serial interface or by a conventional nonvolatile memory programmer. By combining an enhanced RISC 8-bit CPU with In-System Programmable Flash on a monolithic chip, the Atmel AT90S1200 is a powerful microcontroller that provides a highly flexible and cost-effective solution to many embedded control applications.

The AT90S1200 AVR is supported with a full suite of program and system development tools including: macro assemblers, program debugger/simulators, in-circuit emulators, and evaluation kits.

Pin Descriptions

VCC Supply voltage pin.

GND Ground pin.

Port B (PB7..PB0) Port B is an 8-bit bi-directional I/O port. Port pins can provide internal pull-up resistors (selected for each bit). PB0 and PB1 also serve as the positive input (AIN0) and the negative input (AIN1), respectively, of the On-chip Analog Comparator. The Port B output buffers can sink 20 mA and thus drive LED displays directly. When pins PB0 to PB7 are used as inputs and are externally pulled low, they will source current if the internal pull-up resistors are activated. The Port B pins are tri-stated when a reset condition becomes active, even if the clock is not active.

Port B also serves the functions of various special features of the AT90S1200 as listed on page 30.

Port D (PD6..PD0) Port D has seven bi-directional I/O pins with internal pull-up resistors, PD6..PD0. The Port D output buffers can sink 20 mA. As inputs, Port D pins that are externally pulled low will source current if the pull-up resistors are activated. The Port D pins are tri-stated when a reset condition becomes active, even if the clock is not active.

Port D also serves the functions of various special features of the AT90S1200 as listed on page 34.

$\overline{\text{RESET}}$ Reset input. A low level on this pin for more than 50 ns will generate a reset, even if the clock is not running. Shorter pulses are not guaranteed to generate a reset.

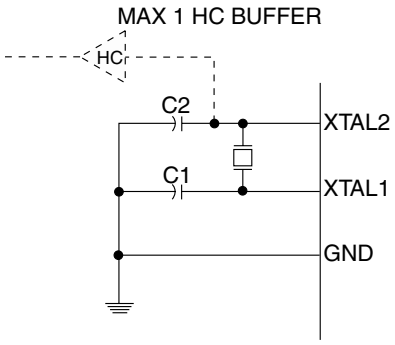
XTAL1 Input to the inverting oscillator amplifier and input to the internal clock operating circuit.

XTAL2 Output from the inverting oscillator amplifier.

Crystal Oscillator

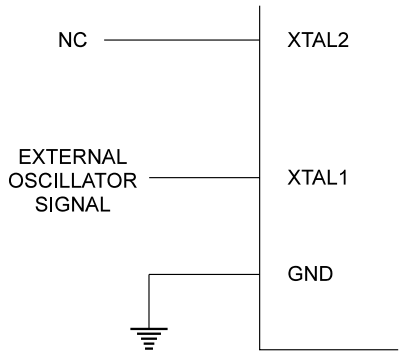
XTAL1 and XTAL2 are input and output, respectively, of an inverting amplifier which can be configured for use as an On-chip Oscillator, as shown in Figure 2. Either a quartz crystal or a ceramic resonator may be used. To drive the device from an external clock source, XTAL2 should be left unconnected while XTAL1 is driven as shown in Figure 3.

Figure 2. Oscillator Connections



Note: When using the MCU Oscillator as a clock for an external device, an HC buffer should be connected as indicated in the figure.

Figure 3. External Clock Drive Configuration



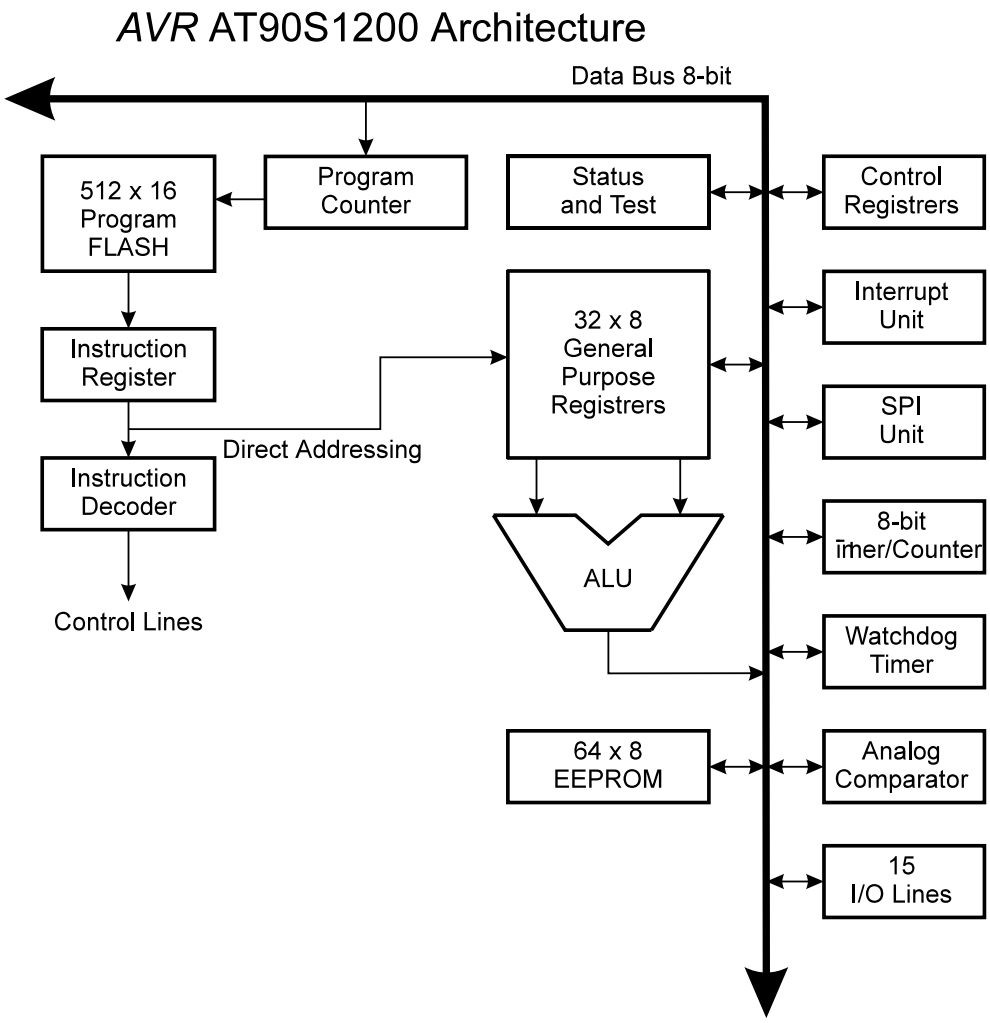
On-chip RC Oscillator

An On-chip RC Oscillator running at a fixed frequency of 1 MHz can be selected as the MCU clock source. If enabled, the AT90S1200 can operate with no external components. A control bit (RCEN) in the Flash Memory selects the On-chip RC Oscillator as the clock source when programmed (“0”). The AT90S1200 is normally shipped with this bit unprogrammed (“1”). Parts with this bit programmed can be ordered as AT90S1200A. The RCEN-bit can be changed by parallel programming only. When using the On-chip RC Oscillator for Serial Program downloading, the RCEN bit must be programmed in Parallel Programming mode first.

Architectural Overview

The fast-access register file concept contains 32 x 8-bit general purpose working registers with a single clock cycle access time. This means that during one single clock cycle, one ALU (Arithmetic Logic Unit) operation is executed. Two operands are output from the register file, the operation is executed, and the result is stored back in the register file – in one clock cycle.

Figure 4. The AT90S1200 AVR RISC Architecture



The ALU supports arithmetic and logic functions between registers or between a constant and a register. Single register operations are also executed in the ALU. Figure 4 shows the AT90S1200 AVR RISC microcontroller architecture. The AVR uses a Harvard architecture concept – with separate memories and buses for program and data memories. The program memory is accessed with a 2-stage pipeline. While one instruction is being executed, the next instruction is pre-fetched from the program memory. This concept enables instructions to be executed in every clock cycle. The program memory is In-System Programmable Flash memory.

With the relative jump and relative call instructions, the whole 512 address space is directly accessed. All AVR instructions have a single 16-bit word format, meaning that every program memory address contains a single 16-bit instruction.

During interrupts and subroutine calls, the return address Program Counter (PC) is stored on the stack. The stack is a 3-level-deep hardware stack dedicated for subroutines and interrupts.

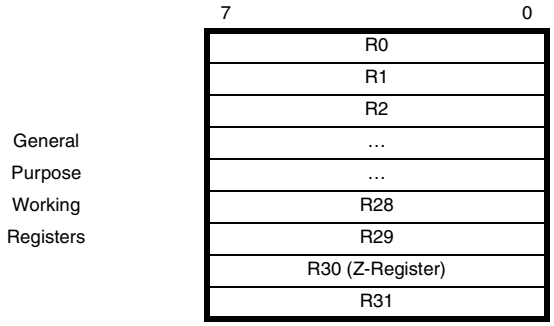
The I/O memory space contains 64 addresses for CPU peripheral functions such as Control Registers, Timer/Counters, A/D Converters and other I/O functions. The memory spaces in the AVR architecture are all linear and regular memory maps.

A flexible interrupt module has its control registers in the I/O space with an additional global interrupt enable bit in the status register. All the different interrupts have a separate interrupt vector in the interrupt vector table at the beginning of the program memory. The different interrupts have priority in accordance with their interrupt vector position. The lower the interrupt vector address, the higher the priority.

General Purpose Register File

Figure 5 shows the structure of the 32 general purpose registers in the CPU.

Figure 5. AVR CPU General Purpose Working Registers



All the register operating instructions in the instruction set have direct and single cycle access to all registers. The only exception is the five constant arithmetic and logic instructions SBCI, SUBI, CPI, ANDI, ORI between a constant and a register and the LDI instruction for load immediate constant data. These instructions apply to the second half of the registers in the register file (R16..R31). The general SBC, SUB, CP, AND, OR and all other operations between two registers or on a single register apply to the entire register file.

Register 30 also serves as an 8-bit pointer for indirect address of the register file.

ALU – Arithmetic Logic Unit

The high-performance AVR ALU operates in direct connection with all the 32 general purpose working registers. Within a single clock cycle, ALU operations between registers in the register file are executed. The ALU operations are divided into three main categories – arithmetic, logic and bit-functions.

In-System Programmable Flash Program Memory

The AT90S1200 contains 1K bytes On-chip In-System Programmable Flash memory for program storage. Since all instructions are single 16-bit words, the Flash is organized as 512 x 16. The Flash memory has an endurance of at least 1000 write/erase cycles.

The AT90S1200 Program Counter is 9 bits wide, thus addressing the 512 words Flash program memory.

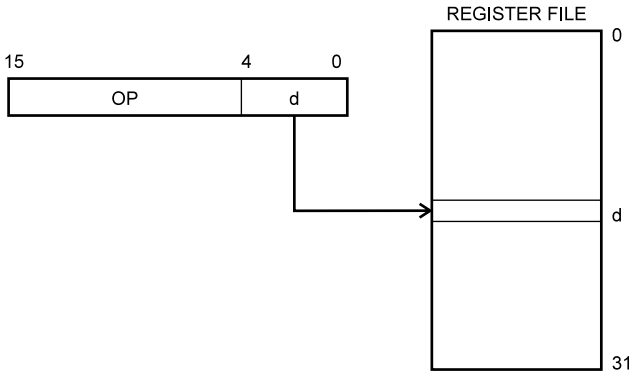
See page 37 for a detailed description on Flash data downloading.

Program and Data Addressing Modes

The AT90S1200 AVR RISC Microcontroller supports powerful and efficient addressing modes. This section describes the different addressing modes supported in the AT90S1200. In the figures, OP means the operation code part of the instruction word. To simplify, not all figures show the exact location of the addressing bits.

Register Direct, Single Register Rd

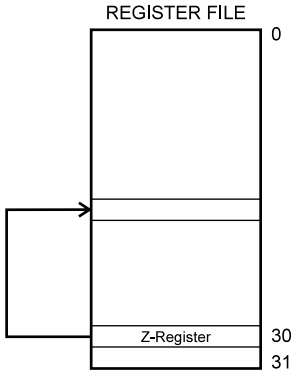
Figure 6. Direct Single Register Addressing



The operand is contained in register d (Rd).

Register Indirect

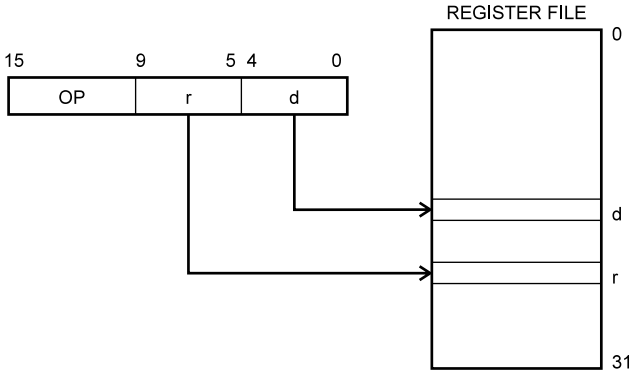
Figure 7. Indirect Register Addressing



The register accessed is the one pointed to by the Z-register (R30).

Register Direct, Two Registers Rd and Rr

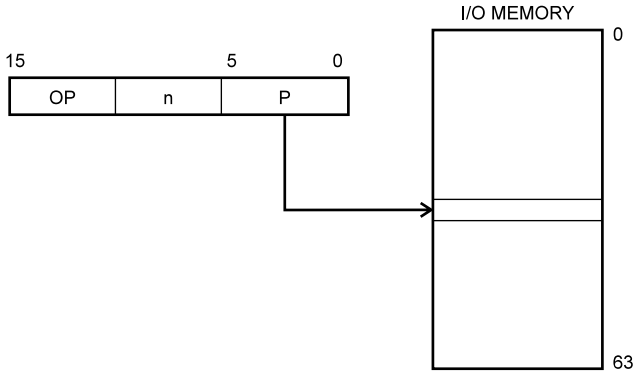
Figure 8. Direct Register Addressing, Two Registers



Operands are contained in register r (Rr) and d (Rd). The result is stored in register d (Rd).

I/O Direct

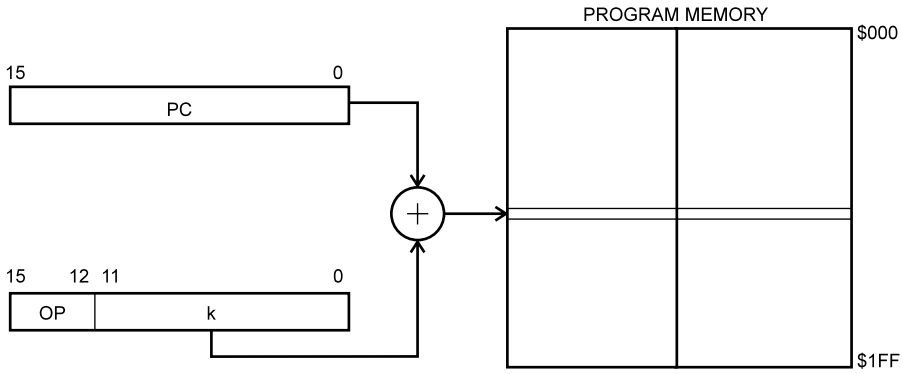
Figure 9. I/O Direct Addressing



Operand address is contained in 6 bits of the instruction word. n is the destination or source register address.

Relative Program Addressing, RJMP and RCALL

Figure 10. Relative Program Memory Addressing



Program execution continues at address PC + k + 1. The relative address k is -2048 to 2047.

Subroutine and Interrupt Hardware Stack

The AT90S1200 uses a 3 level deep hardware stack for subroutines and interrupts. The hardware stack is 9 bits wide and stores the Program Counter (PC) return address while subroutines and interrupts are executed.

RCALL instructions and interrupts push the PC return address onto stack level 0, and the data in the other stack levels 1 - 2 are pushed one level deeper in the stack. When a RET or RETI instruction is executed the returning PC is fetched from stack level 0, and the data in the other stack levels 1 - 2 are popped one level in the stack.

If more than three subsequent subroutine calls or interrupts are executed, the first values written to the stack are overwritten.

EEPROM Data Memory

The AT90S1200 contains 64 bytes of data EEPROM memory. It is organized as a separate data space, in which single bytes can be read and written. The EEPROM has an endurance of at least 100,000 write/erase cycles. The access between the EEPROM and the CPU is described on page 25 specifying the EEPROM address register, the EEPROM data register, and the EEPROM control register. For the SPI data downloading, see page 44 for a detailed description.

Instruction Execution Timing

This section describes the general access timing concepts for instruction execution and internal memory access.

The AVR CPU is driven by the System Clock Ø, directly generated from the external clock crystal for the chip. No internal clock division is used.

Figure 11 shows the parallel instruction fetches and instruction executions enabled by the Harvard architecture and the fast-access register file concept. This is the basic pipelining concept to obtain up to 1 MIPS per MHz with the corresponding unique results for functions per cost, functions per clocks, and functions per power-unit.

Figure 11. The Parallel Instruction Fetches and Instruction Executions

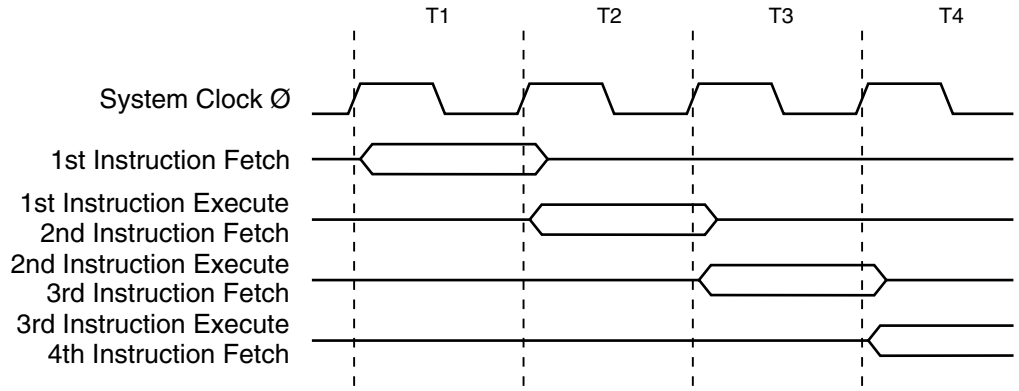
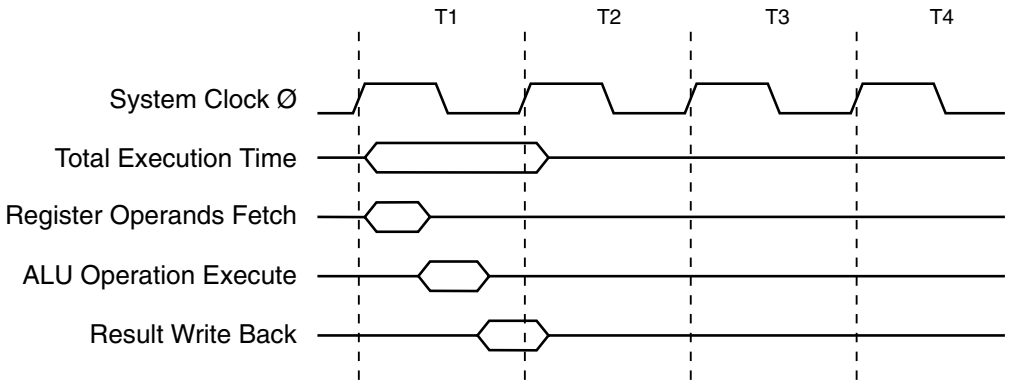


Figure 12 shows the internal timing concept for the register file. In a single clock cycle an ALU operation using two register operands is executed, and the result is stored back to the destination register.

Figure 12. Single-cycle ALU Operation



I/O Memory

The I/O space definition of the AT90S1200 is shown in the following table.

Table 1. The AT90S1200 I/O Space

Address Hex	Name	Function
$3F	SREG	Status REGister
$3B	GIMSK	General Interrupt MaSK register
$39	TIMSK	Timer/Counter Interrupt MaSK register
$38	TIFR	Timer/Counter Interrupt Flag register
$35	MCUCR	MCU general Control Register
$33	TCCR0	Timer/Counter0 Control Register
$32	TCNT0	Timer/Counter0 (8-bit)
$21	WDTCR	Watchdog Timer Control Register
$1E	EEAR	EEPROM Address Register
$1D	EEDR	EEPROM Data Register
$1C	EECR	EEPROM Control Register
$18	PORTB	Data Register, Port B
$17	DDRB	Data Direction Register, Port B
$16	PINB	Input Pins, Port B
$12	PORTD	Data Register, Port D
$11	DDRD	Data Direction Register, Port D
$10	PIND	Input Pins, Port D
$08	ACSR	Analog Comparator Control and Status Register

Note: Reserved and unused locations are not shown in the table.

All AT90S1200 I/Os and peripherals are placed in the I/O space. The different I/O locations are accessed by the IN and OUT instructions transferring data between the 32 general purpose working registers and the I/O space. I/O registers within the address range $00 - $1F are directly bit-accessible using the SBI and CBI instructions. In these registers, the value of single bits can be checked by using the SBIS and SBIC instructions. Refer to the instruction set chapter for more details.

For compatibility with future devices, reserved bits should be written to zero if accessed. Reserved I/O memory addresses should never be written.

Some of the status flags are cleared by writing a logical one to them. Note that the CBI and SBI instructions will operate on all bits in the I/O register, writing a one back into any flag read as set, thus clearing the flag. The CBI and SBI instructions work with registers $00 to $1F only.

The different I/O and peripherals control registers are explained in the following sections.

Status Register – SREG

The AVR status register (SREG) at I/O space location $3F is defined as:

Bit	7	6	5	4	3	2	1	0	
$3F	I	T	H	S	V	N	Z	C	SREG
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	
Initial Value	0	0	0	0	0	0	0	0	

- **Bit 7 – I: Global Interrupt Enable**

The global interrupt enable bit must be set (one) for the interrupts to be enabled. The individual interrupt enable control is then performed in separate control registers. If the global interrupt enable bit is cleared (zero), none of the interrupts are enabled independent of the individual interrupt enable settings. The I-bit is cleared by hardware after an interrupt has occurred, and is set by the RETI instruction to enable subsequent interrupts.

- **Bit 6 – T: Bit Copy Storage**

The bit copy instructions BLD (Bit LoaD) and BST (Bit STore) use the T-bit as source and destination for the operated bit. A bit from a register in the register file can be copied into T by the BST instruction, and a bit in T can be copied into a bit in a register in the register file by the BLD instruction.

- **Bit 5 – H: Half-carry Flag**

The half-carry flag H indicates a half carry in some arithmetic operations. See the Instruction Set description for detailed information.

- **Bit 4 – S: Sign Bit, S = N⊕V**

The S-bit is always an exclusive or between the negative flag N and the two's complement overflow flag V. See the Instruction Set description for detailed information.

- **Bit 3 – V: Two's Complement Overflow Flag**

The two's complement overflow flag V supports two's complement arithmetics. See the Instruction Set description for detailed information.

- **Bit 2 – N: Negative Flag**

The negative flag N indicates a negative result after the different arithmetic and logic operations. See the Instruction Set description for detailed information.

- **Bit 1 – Z: Zero Flag**

The zero flag Z indicates a zero result after the different arithmetic and logic operations. See the Instruction Set description for detailed information.

- **Bit 0 – C: Carry Flag**

The carry flag C indicates a carry in an arithmetic or logic operation. See the Instruction Set description for detailed information.

Note that the status register is not automatically stored when entering an interrupt routine and restored when returning from an interrupt routine. This must be handled by software.

Reset and Interrupt Handling

The AT90S1200 provides three different interrupt sources. These interrupts and the separate reset vector, each have a separate program vector in the program memory space. All the interrupts are assigned individual enable bits that must be set (one) together with the I-bit in the Status Register in order to enable the interrupt.

The lowest addresses in the program memory space are automatically defined as the Reset and Interrupt vectors. The complete list of vectors is shown in Table 2. The list also determines the priority levels of the different interrupts. The lower the address the higher is the priority level. RESET has the highest priority, and next is INT0 (the External Interrupt Request 0), etc.

Table 2. Reset and Interrupt Vectors

Vector No.	Program Address	Source	Interrupt Definition
1	$000	RESET	Hardware Pin, Power-on Reset and Watchdog Reset
2	$001	INT0	External Interrupt Request 0
4	$002	TIMER0, OVF0	Timer/Counter0 Overflow
5	$003	ANA_COMP	Analog Comparator

The most typical and general program setup for the Reset and Interrupt Vector Addresses are:

```
Address   Labels   Code                 Comments
$000               rjmp    RESET        ; Reset Handler
$001               rjmp    EXT_INT0     ; IRQ0 Handler
$002               rjmp    TIM0_OVF     ; Timer0 Overflow Handler
$003               rjmp    ANA_COMP     ; Analog Comparator Handler
;
$004      MAIN:    <instr> xxx          ; Main program start
  …         …        …     …
```

Reset Sources

The AT90S1200 has three sources of reset:

- Power-on Reset. The MCU is reset when the supply voltage is below the power-on Reset threshold (V_{POT}).
- External Reset. The MCU is reset when a low level is present on the $\overline{\text{RESET}}$ pin for more than 50 ns.
- Watchdog Reset. The MCU is reset when the Watchdog Timer period expires and the Watchdog is enabled.

During Reset, all I/O registers are then set to their initial values, and the program starts execution from address $000. The instruction placed in address $000 must be an RJMP (relative jump) instruction to the reset handling routine. If the program never enables an interrupt source, the interrupt vectors are not used, and regular program code can be placed at these locations. The circuit diagram in Figure 13 shows the reset logic. Table 3 defines the timing and electrical parameters of the reset circuitry. Note that Power-on Reset timing is clocked by the internal RC Oscillator. Refer to characterization data for RC Oscillator frequency at other V_{CC} voltages.

Figure 13. Reset Logic

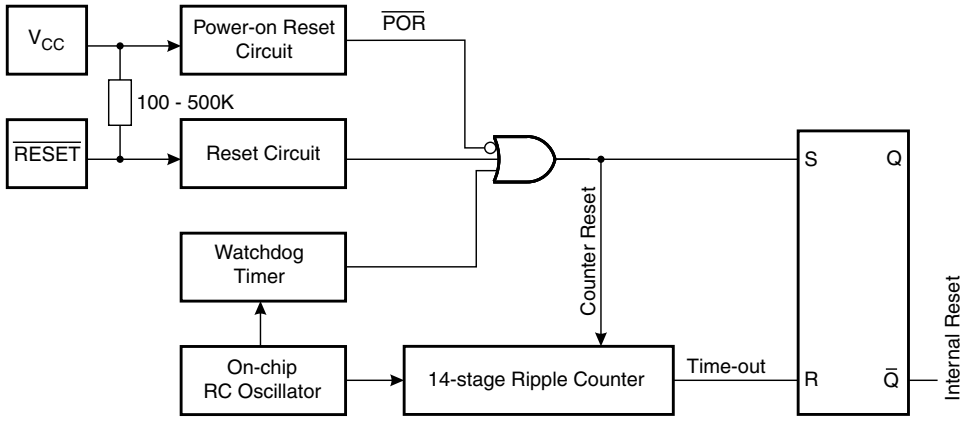


Table 3. Reset Characteristics (V_{CC} = 5.0V)

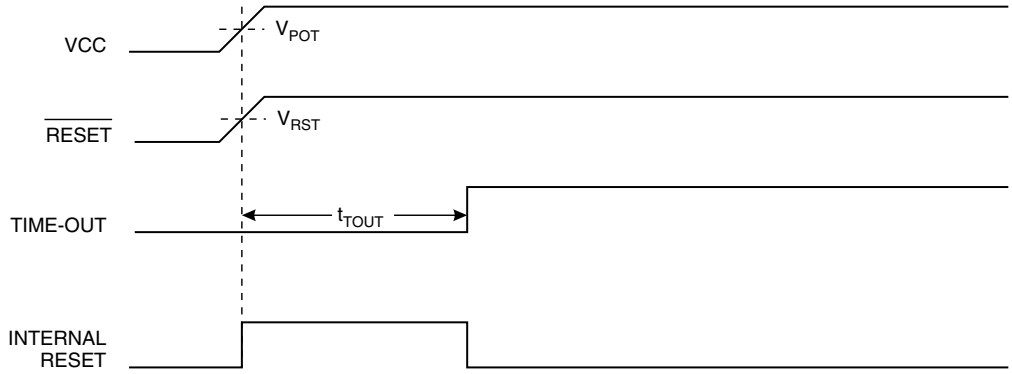
Symbol	Parameter	Min	Typ	Max	Units
V_{POT}(1)	Power-on Reset Threshold Voltage (rising)	0.8	1.2	1.6	V
	Power-on Reset Threshold Voltage (falling)	0.2	0.4	0.6	V
V_{RST}	Pin Threshold Voltage	–	–	0.85 V_{CC}	V
t_{POR}	Power-on Reset Period	2.0	3.0	4.0	ms
t_{TOUT}	Reset Delay Time-out Period (The Time-out period equals 16K WDT cycles. See “Typical Characteristics” on page 51. for typical WDT frequency at different voltages).	11.0	16.0	21.0	ms

Note: 1. The Power-on Reset will not work unless the supply voltage has been below V_{POT} (falling).

Power-on Reset

A Power-on Reset (POR) circuit ensures that the device is reset from power-on. As shown in Figure 13, an internal timer clocked from the Watchdog timer oscillator prevents the MCU from starting until after a certain period after V_{CC} has reached the Power-on Threshold voltage (V_{POT}), regardless of the V_{CC} rise time (see Figure 14).

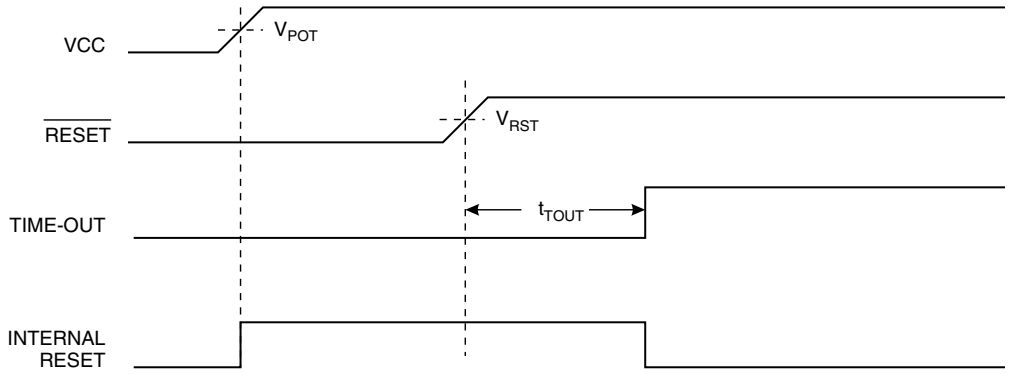
Figure 14. MCU Start-up, $\overline{RESET}$ Tied to V_{CC}.



If the built-in start-up delay is sufficient, $\overline{RESET}$ can be connected to V_{CC} directly or via an external pull-up resistor. By holding the $\overline{RESET}$ pin low for a period after V_{CC} has

been applied, the Power-on Reset period can be extended. Refer to Figure 15 for a timing example on this.

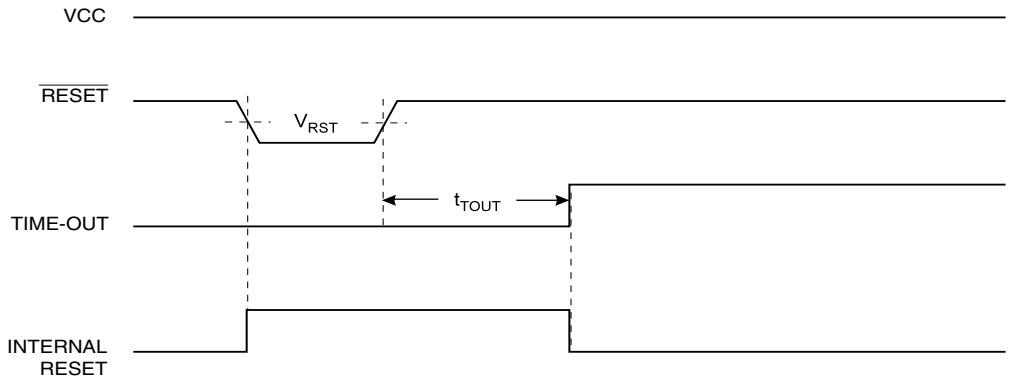
Figure 15. MCU Start-up, $\overline{\text{RESET}}$ Controlled Externally



External Reset

An External Reset is generated by a low level on the $\overline{\text{RESET}}$ pin. Reset pulses longer than 50 ns will generate a reset, even if the clock is not running. Shorter pulses are not guaranteed to generate a reset. When the applied signal reaches the Reset Threshold Voltage (V_{RST}) on its positive edge, the delay timer starts the MCU after the Time-out period t_{TOUT} has expired.

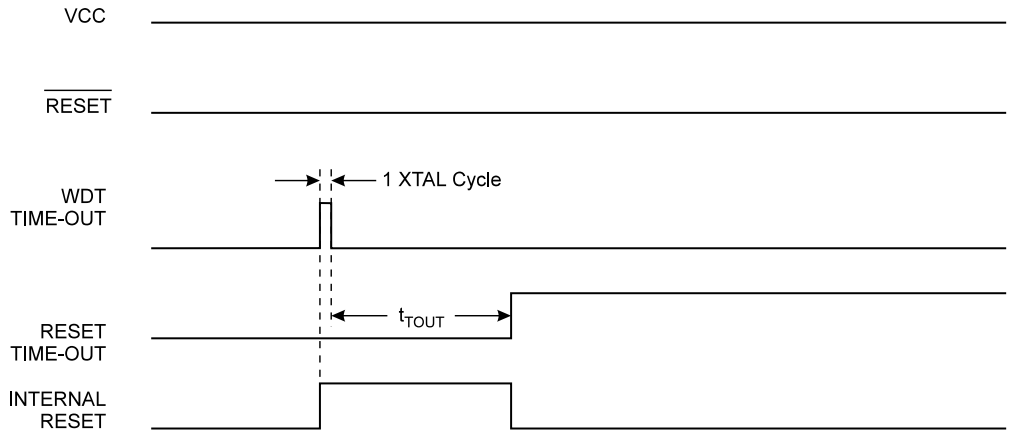
Figure 16. External Reset during Operation



Watchdog Reset

When the Watchdog times out, it will generate a short reset pulse of 1 XTAL cycle duration. On the falling edge of this pulse, the delay timer starts counting the Time-out period t_{TOUT}. Refer to page 23 for details on operation of the Watchdog.

Figure 17. Watchdog Reset during Operation



Interrupt Handling

The AT90S1200 has two Interrupt Mask Control Registers: the GIMSK (General Interrupt Mask Register) at I/O space address $3B and the TIMSK (Timer/Counter Interrupt Mask Register) at I/O address $39.

When an interrupt occurs, the Global Interrupt Enable I-bit is cleared (zero) and all interrupts are disabled. The user software can set (one) the I-bit to enable interrupts. The I-bit is set (one) when a Return from Interrupt instruction (RETI) is executed.

When the Program Counter is vectored to the actual interrupt vector in order to execute the interrupt handling routine, hardware clears the corresponding flag that generated the interrupt. Some of the interrupt flags can also be cleared by writing a logic one to the flag bit position(s) to be cleared.

If an interrupt condition occurs when the corresponding interrupt enable bit is cleared (zero), the interrupt flag will be set and remembered until the interrupt is enabled, or the flag is cleared by software.

If one or more interrupt conditions occur when the global interrupt enable bit is cleared (zero), the corresponding interrupt flag(s) will be set and remembered until the global interrupt enable bit is set (one), and will be executed by order of priority.

Note that external level interrupt does not have a flag, and will only be remembered for as long as the interrupt condition is active.

Note that the Status Register is not automatically stored when entering an interrupt routine and restored when returning from an interrupt routine. This must be handled by software.

General Interrupt Mask Register – GIMSK

Bit	7	6	5	4	3	2	1	0	
$3B	-	INT0	-	-	-	-	-	-	GIMSK
Read/Write	R	R/W	R	R	R	R	R	R	
Initial Value	0	0	0	0	0	0	0	0	

- **Bit 7 – Res: Reserved Bit**

This bit is a reserved bit in the AT90S1200 and always reads as zero.

- **Bit 6 – INT0: External Interrupt Request 0 Enable**

When the INT0 bit is set (one) and the I-bit in the Status Register (SREG) is set (one), the external pin interrupt is enabled. The Interrupt Sense Control0 bit 1/0 (ISC01 and ISC00) in the MCU general Control Register (MCUCR) defines whether the external interrupt is activated on rising or falling edge of the INT0 pin or low level sensed. INT0 can be activated even if the pin is configured as an output. See also page 17.

- **Bits 5..0 – Res: Reserved Bits**

These bits are reserved bits in the AT90S1200 and always read as zero.

Timer/Counter Interrupt Mask Register – TIMSK

Bit	7	6	5	4	3	2	1	0	
$39	-	-	-	-	-	-	TOIE0	-	TIMSK
Read/Write	R	R	R	R	R	R	R/W	R	
Initial Value	0	0	0	0	0	0	0	0	

- **Bits 7..2 – Res: Reserved Bits**

These bits are reserved bits in the AT90S1200 and always read as zero.

- **Bit 1 – TOIE0: Timer/Counter0 Overflow Interrupt Enable**

When the TOIE0 bit is set (one) and the I-bit in the Status Register is set (one), the Timer/Counter0 Overflow interrupt is enabled. The corresponding interrupt (at vector $002) is executed if an overflow in Timer/Counter0 occurs, i.e., when the TOV0 bit is set in the Timer/Counter Interrupt Flag Register (TIFR).

- **Bit 0 – Res: Reserved Bit**

This bit is a reserved bit in the AT90S1200 and always reads as zero.

Timer/Counter Interrupt FLAG Register – TIFR

Bit	7	6	5	4	3	2	1	0	
$38	-	-	-	-	-	-	TOV0	-	TIFR
Read/Write	R	R	R	R	R	R	R/W	R	
Initial Value	0	0	0	0	0	0	0	0	

- **Bits 7..2 – Res: Reserved Bits**

These bits are reserved bits in the AT90S1200 and always read as zero.

- **Bit 1 – TOV0: Timer/Counter0 Overflow Flag**

The bit TOV0 is set (one) when an overflow occurs in Timer/Counter0. TOV0 is cleared by hardware when executing the corresponding interrupt handling vector. Alternatively, TOV0 is cleared by writing a logic one to the flag. When the SREG I-bit, and TOIE0 (Timer/Counter0 Overflow Interrupt Enable), and TOV0 are set (one), the Timer/Counter0 Overflow interrupt is executed.

- **Bit 0 – Res: Reserved Bit**

This bit is a reserved bit in the AT90S1200 and always reads as zero.

External Interrupts

The External Interrupt is triggered by the INT0 pin. The interrupt can trigger on rising edge, falling edge or low level. This is set up as described in the specification for the MCU Control Register (MCUCR). When INT0 is level triggered, the interrupt is pending as long as INT0 is held low.

The interrupt is triggered even if INT0 is configured as an output. This provides a way to generate a software interrupt.

The interrupt flag can not be directly accessed by the user. If an external edge-triggered interrupt is suspected to be pending, the flag can be cleared as follows.

1. Disable the External Interrupt by clearing the INT0 flag in GIMSK.
2. Select level triggered interrupt.
3. Select desired interrupt edge.
4. Re-enable the external interrupt by setting INT0 in GIMSK.

Interrupt Response Time

The interrupt execution response for all the enabled AVR interrupts is four clock cycles minimum. Four clock cycles after the interrupt flag has been set, the program vector address for the actual interrupt handling routine is executed. During this 4-clock-cycle period, the Program Counter (9 bits) is pushed onto the Stack. The vector is normally a relative jump to the interrupt routine, and this jump takes two clock cycles. If an interrupt occurs during execution of a multi-cycle instruction, this instruction is completed before the interrupt is served.

A return from an interrupt handling routine takes four clock cycles. During these four clock cycles, the Program Counter (9 bits) is popped back from the Stack and the I-flag in SREG is set. When the AVR exits from an interrupt, it will always return to the main program and execute one more instruction before any pending interrupt is served.

Note that the Subroutine and Interrupt Stack is a 3-level true hardware stack, and if more than three nested subroutines and interrupts are executed, only the most recent three return addresses are stored.

MCU Control Register – MCUCR

The MCU Control Register contains general microcontroller control bits for general MCU control functions.

Bit	7	6	5	4	3	2	1	0	
$35	–	–	SE	SM	–	–	ISC01	ISC00	MCUCR
Read/Write	R	R	R/W	R/W	R	R	R/W	R/W	
Initial Value	0	0	0	0	0	0	0	0	

- **Bits 7, 6 – Res: Reserved Bits**

These bits are reserved bits in the AT90S1200 and always read as zero.

- **Bit 5 – SE: Sleep Enable**

The SE bit must be set (one) to make the MCU enter the Sleep mode when the SLEEP instruction is executed. To avoid the MCU entering the Sleep mode unless it is the programmers purpose, it is recommended to set the Sleep Enable SE bit just before the execution of the SLEEP instruction.

- **Bit 4 – SM: Sleep Mode**

This bit selects between the two available sleep modes. When SM is cleared (zero), Idle mode is selected as sleep mode. When SM is set (one), Power-down mode is selected as sleep mode. For details, refer to the paragraph “Sleep Modes” on the following page.

- **Bits 3, 2 – Res: Reserved Bits**

These bits are reserved bits in the AT90S1200 and always read as zero.

- **Bits 1, 0 – ISC01, ISC00: Interrupt Sense Control 0 Bit 1 and Bit 0**

The External Interrupt 0 is activated by the external pin INT0 if the SREG I-flag and the corresponding interrupt mask in the GIMSK register is set. The level and edges on the external INT0 pin that activate the interrupt are defined in Table 4.

Table 4. Interrupt 0 Sense Control

ISC01	ISC00	Description
0	0	The low level of INT0 generates an interrupt request.
0	1	Reserved
1	0	The falling edge of INT0 generates an interrupt request.
1	1	The rising edge of INT0 generates an interrupt request.

The value on the INT0 pin is sampled before detecting edges. If edge interrupt is selected, pulses with a duration longer than one CPU clock period will generate an interrupt. Shorter pulses are not guaranteed to generate an interrupt. If low level interrupt is selected, the low level must be held until the completion of the currently executing instruction to generate an interrupt. If enabled, a level triggered interrupt will generate an interrupt request as long as the pin is held low.

Sleep Modes

To enter the sleep modes, the SE bit in MCUCR must be set (one) and a SLEEP instruction must be executed. If an enabled interrupt occurs while the MCU is in a sleep mode, the MCU awakes, executes the interrupt routine, and resumes execution from the instruction following SLEEP. The contents of the register file and the I/O memory are unaltered. If a Reset occurs during sleep mode, the MCU wakes up and executes from the Reset Vector.

Idle Mode

When the SM bit is cleared (zero), the SLEEP instruction makes the MCU enter the Idle mode, stopping the CPU but allowing Timer/Counters, Watchdog and the interrupt system to continue operating. This enables the MCU to wake up from external triggered interrupts as well as internal ones like Timer Overflow interrupt and Watchdog Reset. If wakeup from the Analog Comparator interrupt is not required, the Analog Comparator can be powered down by setting the ACD-bit in the Analog Comparator Control and Status Register (ACSR). This will reduce power consumption in Idle mode. When the MCU wakes up from Idle mode, the CPU starts program execution immediately.

Power-down Mode

When the SM bit is set (one), the SLEEP instruction makes the MCU enter Power-down mode. In this mode, the External Oscillator is stopped while the External Interrupts and the Watchdog (if enabled) continue operating. Only an External Reset, a Watchdog Reset (if enabled), an external level interrupt on INT0 can wake up the MCU.

Note that when a level triggered interrupt is used for wake-up from Power-down, the low level must be held for a time longer than the reset delay time-out period t_{TOUT}. Otherwise, the device will not wake up.

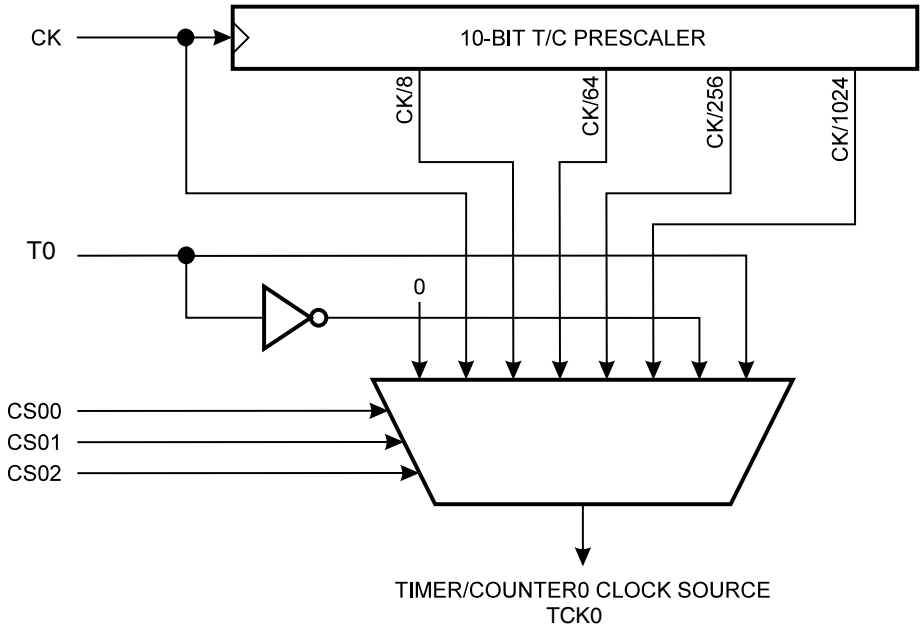
Timer/Counter0

The AT90S1200 provides one general purpose 8-bit Timer/Counter. The Timer/Counter0 gets the prescaled clock from the 10-bit prescaling timer. The Timer/Counter0 can either be used as a Timer with an internal clock time base or as a Counter with an external pin connection, which triggers the counting.

Timer/Counter0 Prescaler

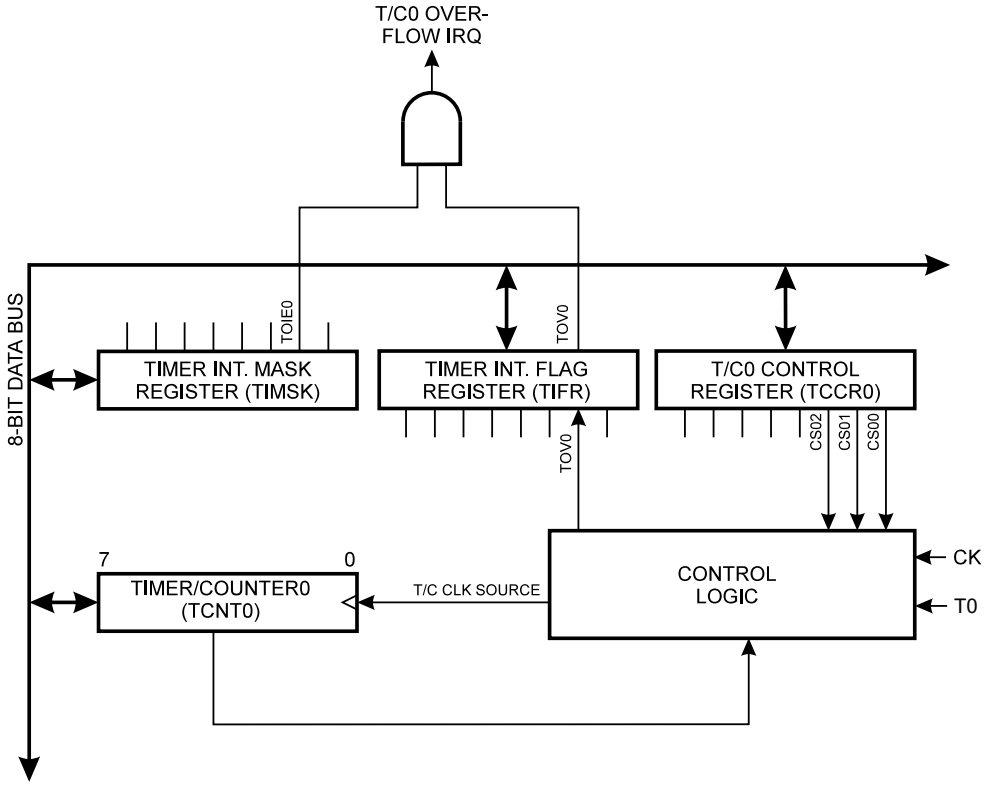
Figure 18 shows the general Timer/Counter0 prescaler.

Figure 18. Timer/Counter0 Prescaler



The four different prescaled selections are: CK/8, CK/64, CK/256, and CK/1024 where CK is the Oscillator Clock. For the Timer/Counter0, added selections as CK, external clock source and stop, can be selected as clock sources. Figure 19 shows the block diagram for Timer/Counter0.

Figure 19. Timer/Counter0 Block Diagram



The 8-bit Timer/Counter0 can select clock source from CK, prescaled CK or an external pin. In addition it can be stopped as described in the specification for the Timer/Counter0 Control Register (TCCR0). The overflow status flag is found in the Timer/Counter Interrupt Flag Register (TIFR). Control signals are found in the Timer/Counter0 Control Register (TCCR0). The interrupt enable/disable settings for Timer/Counter0 are found in the Timer/Counter Interrupt Mask Register (TIMSK).

When Timer/Counter0 is externally clocked, the external signal is synchronized with the oscillator frequency of the CPU. To assure proper sampling of the external clock, the minimum time between two external clock transitions must be at least one internal CPU clock period. The external clock signal is sampled on the rising edge of the internal CPU clock.

The 8-bit Timer/Counter0 features both a high-resolution and a high-accuracy usage with the lower prescaling opportunities. Similarly, the high prescaling opportunities make the Timer/Counter0 useful for lower speed functions or exact timing functions with infrequent actions.

Timer/Counter0 Control Register – TCCR0

Bit	7	6	5	4	3	2	1	0	
$33	-	-	-	-	-	CS02	CS01	CS00	TCCR0
Read/Write	R	R	R	R	R	R/W	R/W	R/W	
Initial Value	0	0	0	0	0	0	0	0	

- **Bits 7..3 – Res: Reserved Bits**

These bits are reserved bits in the AT90S1200 and always read as zero.

- **Bits 2, 1, 0 – CS02, CS01, CS00: Clock Select0, Bits 2, 1 and 0**

The Clock Select0 bits 2, 1 and 0 define the prescaling source of Timer/Counter0.

Table 5. Clock 0 Prescale Select

CS02	CS01	CS00	Description
0	0	0	Stop, the Timer/Counter0 is stopped.
0	0	1	CK
0	1	0	CK/8
0	1	1	CK/64
1	0	0	CK/256
1	0	1	CK/1024
1	1	0	External Pin T0, falling edge
1	1	1	External Pin T0, rising edge

The Stop condition provides a Timer Enable/Disable function. The CK down divided modes are scaled directly from the CK Oscillator clock. If the external pin modes are used for Timer/Counter0, transitions on PD4/(T0) will clock the counter even if the pin is configured as an output. This feature can give the user SW control of the counting.

Timer/Counter0 – TCNT0

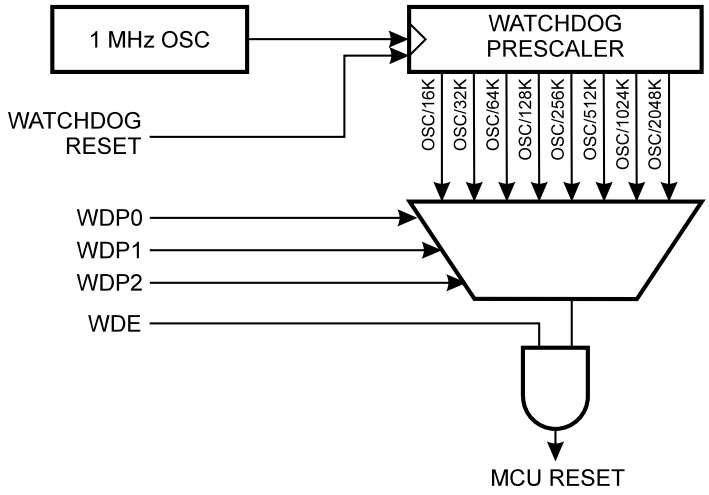
Bit	7	6	5	4	3	2	1	0	
$32	MSB							LSB	TCNT0
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	
Initial Value	0	0	0	0	0	0	0	0	

The Timer/Counter0 is realized as an up-counter with read and write access. If the Timer/Counter0 is written and a clock source is present, the Timer/Counter0 continues counting in the timer clock cycle following the write operation.

Watchdog Timer

The Watchdog Timer is clocked from a separate On-chip Oscillator that runs at 1 MHz. This is the typical value at V_{CC} = 5V. See characterization data for typical values at other V_{CC} levels. By controlling the Watchdog Timer prescaler, the Watchdog Reset interval can be adjusted, see Table 6 for a detailed description. The WDR (Watchdog Reset) instruction resets the Watchdog Timer. Eight different clock cycle periods can be selected to determine the maximum period between two WDR instructions to prevent the Watchdog Timer from resetting the MCU. If the reset period expires without another WDR instruction, the AT90S1200 resets and executes from the Reset Vector. For timing details on the Watchdog Reset, refer to page 14.

Figure 20. Watchdog Timer



Watchdog Timer Control Register – WDTCR

Bit	7	6	5	4	3	2	1	0	
$21	–	–	–	–	WDE	WDP2	WDP1	WDP0	WDTCR
Read/Write	R	R	R	R	R/W	R/W	R/W	R/W	
Initial Value	0	0	0	0	0	0	0	0	

- **Bits 7..4 – Res: Reserved Bits**

These bits are reserved bits in the AT90S1200 and will always read as zero.

- **Bit 3 – WDE: Watchdog Enable**

When the WDE is set (one) the Watchdog Timer is enabled, and if the WDE is cleared (zero) the Watchdog Timer function is disabled.

- **Bits 2..0 – WDP2..0: Watchdog Timer Prescaler 2, 1 and 0**

The WDP2..0 determine the Watchdog Timer prescaling when the Watchdog Timer is enabled. The different prescaling values and their corresponding timeout periods are shown in Table 6.

Table 6. Watchdog Timer Prescale Select

WDP2	WDP1	WDP0	Number of WDT Oscillator Cycles	Typical Time-out at V_{CC} = 3.0V	Typical Time-out at V_{CC} = 5.0V
0	0	0	16K cycles	47 ms	15 ms
0	0	1	32K cycles	94 ms	30 ms
0	1	0	64K cycles	0.19 s	60 ms
0	1	1	128K cycles	0.38 s	0.12 s
1	0	0	256K cycles	0.75 s	0,24 s
1	0	1	512K cycles	1.5 s	0.49 s
1	1	0	1,024K cycles	3.0 s	0.97 s
1	1	1	2,048K cycles	6.0 s	1.9 s

Note: The frequency of the Watchdog Oscillator is voltage dependent as shown in "Typical Characteristics" on page 51.

The WDR (Watchdog Reset) instruction should always be executed before the Watchdog Timer is enabled. This ensures that the reset period will be in accordance with the Watchdog Timer prescale settings. If the Watchdog Timer is enabled without Reset, the Watchdog Timer may not start to count from zero.

To avoid unintentional MCU resets, the Watchdog Timer should be disabled or reset before changing the Watchdog Timer Prescale Select.

EEPROM Read/Write Access

The EEPROM access registers are accessible in the I/O space.

The write access time is in the range of 2.5 - 4 ms, depending on the V_{CC} voltages. A self-timing function, however, lets the user software detect when the next byte can be written. If the user code contains code that writes the EEPROM, some precaution must be taken. In heavily filtered power supplies, V_{CC} is likely to rise or fall slowly on Power-up/down. This causes the device for some period of time to run at a voltage lower than specified as minimum for the clock frequency used. CPU operation under these conditions is likely cause the program counter to perform unintentional jumps and eventually execute the EEPROM write code. To secure EEPROM integrity, the user is advised to use an external under-voltage reset circuit in this case.

In order to prevent unintentional EEPROM writes, a specific write procedure must be followed. Refer to “EEPROM Control Register – EECR” on page 25 for details on this.

When the EEPROM is read or written, the CPU is halted for two clock cycles before the next instruction is executed.

EEPROM Address Register – EEAR

Bit	7	6	5	4	3	2	1	0	
$1E	–	–	EEAR5	EEAR4	EEAR3	EEAR2	EEAR1	EEAR0	EEAR
Read/Write	R	R	R/W	R/W	R/W	R/W	R/W	R/W	
Initial Value	0	0	0	0	0	0	0	0	

- **Bit 7, 6 – Res: Reserved Bits**

These bits are reserved bits in the AT90S1200 and will always read as zero.

- **Bits 5..0 – EEAR5..0: EEPROM Address**

The EEPROM Address Register (EEAR5..0) specifies the EEPROM address in the 64-byte EEPROM space. The EEPROM data bytes are addressed linearly between 0 and 63.

EEPROM Data Register – EEDR

Bit	7	6	5	4	3	2	1	0	
$1D	MSB							LSB	EEDR
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	
Initial Value	0	0	0	0	0	0	0	0	

- **Bits 7..0 – EEDR7..0: EEPROM Data**

For the EEPROM write operation, the EEDR register contains the data to be written to the EEPROM in the address given by the EEAR register. For the EEPROM read operation, the EEDR contains the data read out from the EEPROM at the address given by EEAR.

EEPROM Control Register – EECR

Bit	7	6	5	4	3	2	1	0	
$1C	–	–	–	–	–	–	EEWE	EERE	EECR
Read/Write	R	R	R	R	R	R	R/W	R/W	
Initial Value	0	0	0	0	0	0	0	0	

- **Bits 7..2 – Res: Reserved Bits**

These bits are reserved bits in the AT90S1200 and will always be read as zero.

- **Bit 1 – EEWE: EEPROM Write Enable**

The EEPROM Write Enable Signal (EEWE) is the write strobe to the EEPROM. When address and data are correctly set up, the EEWE bit must be set to write the value into the EEPROM. When the write access time (typically 2.5 ms at V_{CC} = 5V and 4 ms at V_{CC} = 2.7V) has elapsed, the EEWE bit is cleared (zero) by hardware. The user software can poll this bit and wait for a zero before writing the next byte. When EEWE has been set, the CPU is halted for two cycles before the next instruction is executed.

- **Bit 0 – EERE: EEPROM Read Enable**

The EEPROM Read Enable Signal (EERE) is the read strobe to the EEPROM. When the correct address is set up in the EEAR register, the EERE bit must be set. When the EERE bit is cleared (zero) by hardware, requested data is found in the EEDR register. The EEPROM read access takes one instruction and there is no need to poll the EERE bit. When EERE has been set, the CPU is halted for four cycles before the next instruction is executed.

Caution: If an interrupt routine accessing the EEPROM is interrupting another EEPROM access, the EEAR or EEDR register will be modified, causing the interrupted EEPROM access to fail. It is recommended to have the global interrupt flag cleared during EEPROM write operation to avoid these problems.

Prevent EEPROM Corruption

During periods of low V_{CC}, the EEPROM data can be corrupted because the supply voltage is too low for the CPU and the EEPROM to operate properly. These issues are the same as for board-level systems using the EEPROM, and the same design solutions should be applied.

An EEPROM data corruption can be caused by two situations when the voltage is too low. First, a regular write sequence to the EEPROM requires a minimum voltage to operate correctly. Secondly, the CPU itself can execute instructions incorrectly, if the supply voltage for executing instructions is too low.

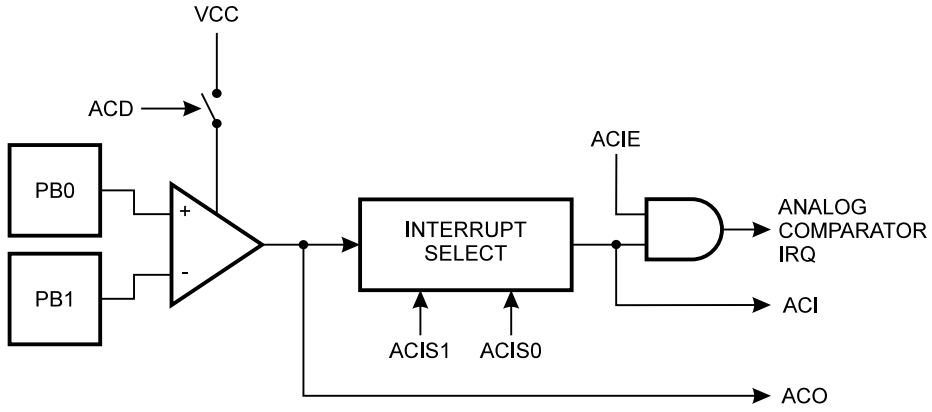
EEPROM data corruption can easily be avoided by following these design recommendations (one is sufficient):

1. Keep the AVR RESET active (low) during periods of insufficient power supply voltage. This is best done by an external low V_{CC} Reset Protection circuit, often referred to as a Brown-out Detector (BOD). Please refer to application note AVR 180 for design considerations regarding power-on reset and low-voltage detection.
2. Keep the AVR core in Power-down Sleep mode during periods of low V_{CC}. This will prevent the CPU from attempting to decode and execute instructions, effectively protecting the EEPROM registers from unintentional writes.
3. Store constants in Flash memory if the ability to change memory contents from software is not required. Flash memory cannot be updated by the CPU, and will not be subject to corruption.

Analog Comparator

The Analog Comparator compares the input values on the positive input PB0 (AIN0) and the negative input PB1 (AIN1). When the voltage on the positive input PB0 (AIN0) is higher than the voltage on the negative input PB1 (AIN1), the Analog Comparator Output (ACO) is set (one). The comparator's output can be set to trigger the Analog Comparator interrupt. The user can select interrupt triggering on comparator output rise, fall or toggle. A block diagram of the comparator and its surrounding logic is shown in Figure 21.

Figure 21. Analog Comparator Block Diagram



Analog Comparator Control and Status Register – ACSR

Bit	7	6	5	4	3	2	1	0	
$08	**ACD**	**–**	**ACO**	**ACI**	**ACIE**	**–**	**ACIS1**	**ACIS0**	**ACSR**
Read/Write	R/W	R	R	R/W	R/W	R	R/W	R/W	
Initial Value	0	0	N/A	0	0	0	0	0	

- **Bit 7 – ACD: Analog Comparator Disable**

When this bit is set (one), the power to the Analog Comparator is switched off. This bit can be set at any time to turn off the analog comparator. This will reduce power consumption in Active and Idle modes. When changing the ACD bit, the Analog Comparator Interrupt must be disabled by clearing the ACIE bit in ACSR. Otherwise, an interrupt can occur when the bit is changed.

- **Bit 6 – Res: Reserved Bit**

This bit is a reserved bit in the AT90S1200 and will always read as zero.

- **Bit 5 – ACO: Analog Comparator Output**

ACO is directly connected to the comparator output.

- **Bit 4 – ACI: Analog Comparator Interrupt Flag**

This bit is set (one) when a comparator output event triggers the interrupt mode defined by ACIS1 and ACIS0. The Analog Comparator Interrupt routine is executed if the ACIE bit is set (one) and the I-bit in SREG is set (one). ACI is cleared by hardware when executing the corresponding interrupt handling vector. Alternatively, ACI is cleared by writing a logic one to the flag. Observe however, that if another bit in this register is modified using the SBI or CBI instruction, ACI will be cleared if it has become set before the operation.

- **Bit 3 – ACIE: Analog Comparator Interrupt Enable**

When the ACIE bit is set (one) and the I-bit in the Status Register is set (one), the Analog Comparator Interrupt is activated. When cleared (zero), the interrupt is disabled.

- **Bit 2 – Res: Reserved Bit**

This bit is a reserved bit in the AT90S1200 and will always read as zero.

- **Bits 1, 0 – ACIS1, ACIS0: Analog Comparator Interrupt Mode Select**

These bits determine which comparator events trigger the Analog Comparator Interrupt. The different settings are shown in Table 7.

Table 7. ACIS1/ACIS0 Settings

ACIS1	ACIS0	Interrupt Mode
0	0	Comparator Interrupt on Output Toggle
0	1	Reserved
1	0	Comparator Interrupt on Falling Output Edge
1	1	Comparator Interrupt on Rising Output Edge

Note: When changing the ACIS1/ACIS0 bits, the Analog Comparator Interrupt must be disabled by clearing its Interrupt Enable bit in the ACSR register. Otherwise, an interrupt can occur when the bits are changed.

I/O Ports

All AVR ports have true Read-Modify-Write functionality when used as general digital I/O ports. This means that the direction of one port pin can be changed without unintentionally changing the direction of any other pin with the SBI and CBI instructions. The same applies for changing drive value (if configured as output) or enabling/disabling of pull-up resistors (if configured as input).

Port B

Port B is an 8-bit bi-directional I/O port.

Three I/O memory address locations are allocated for the Port B, one each for the Data Register – PORTB ($18), Data Direction Register – DDRB ($17), and the Port B Input Pins – PINB ($16). The Port B Input Pins address is read-only, while the Data Register and the Data Direction Register are read/write.

All port pins have individually selectable pull-up resistors. The Port B output buffers can sink 20 mA and thus drive LED displays directly. When pins PB0 to PB7 are used as inputs and are externally pulled low, they will source current if the internal pull-up resistors are activated.

The Port B pins with alternate functions are shown in Table 8.

Table 8. Port B Pin Alternate Functions

Port Pin	Alternate Functions
PB0	AIN0 (Analog Comparator positive input)
PB1	AIN1 (Analog Comparator negative input)
PB5	MOSI (Data Input line for memory downloading)
PB6	MISO (Data Output line for memory uploading)
PB7	SCK (Serial Clock input)

When the pins are used for the alternate function, the DDRB and PORTB register has to be set according to the alternate function description.

Port B Data Register – PORTB

Bit	7	6	5	4	3	2	1	0	
$18	PORTB7	PORTB6	PORTB5	PORTB4	PORTB3	PORTB2	PORTB1	PORTB0	PORTB
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	
Initial Value	0	0	0	0	0	0	0	0	

Port B Data Direction Register – DDRB

Bit	7	6	5	4	3	2	1	0	
$17	DDB7	DDB6	DDB5	DDB4	DDB3	DDB2	DDB1	DDB0	DDRB
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	
Initial Value	0	0	0	0	0	0	0	0	

Port B Input Pin Address – PINB

Bit	7	6	5	4	3	2	1	0	
$16	PINB7	PINB6	PINB5	PINB4	PINB3	PINB2	PINB1	PINB0	PINB
Read/Write	R	R	R	R	R	R	R	R	
Initial Value	N/A	N/A	N/A	N/A	N/A	N/A	N/A	N/A	

The Port B Input Pins address (PINB) is not a register, and this address enables access to the physical value on each Port B pin. When reading PORTB, the Port B Data Latch is read, and when reading PINB, the logical values present on the pins are read.

Port B as General Digital I/O

All eight pins in Port B have equal functionality when used as digital I/O pins.

PBn, General I/O pin: The DDBn bit in the DDRB Register selects the direction of this pin, if DDBn is set (one), PBn is configured as an output pin. If DDBn is cleared (zero), PBn is configured as an input pin. If PORTBn is set (one) and the pin is configured as an input pin, the MOS pull-up resistor is activated. To switch the pull-up resistor off, PORTBn has to be cleared (zero) or the pin has to be configured as an output pin. The Port B pins are tri-stated when a reset condition becomes active, even if the clock is not active.

Table 9. DDBn Effect on Port B Pins

DDBn	PORTBn	I/O	Pull-up	Comment
0	0	Input	No	Tri-state (High-Z)
0	1	Input	Yes	PBn will source current if ext. pulled low.
1	0	Output	No	Push-pull Zero Output
1	1	Output	No	Push-pull One Output

Note: n: 7,6...0, pin number.

Alternate Functions of Port B

The alternate pin functions of Port B are:

- **SCK – Port B, Bit 7**

SCK, Clock Input pin for memory up/downloading.

- **MISO – Port B, Bit 6**

MISO, Data Output pin for memory uploading.

- **MOSI – Port B, Bit 5**

MOSI, Data Input pin for memory downloading.

- **AIN1 – Port B, Bit 1**

AIN1, Analog Comparator Negative Input. When configured as an input (DDB1 is cleared [zero]) and with the internal MOS pull-up resistor switched off (PB1 is cleared [zero]), this pin also serves as the negative input of the On-chip Analog Comparator.

- **AIN0 – Port B, Bit 0**

AIN0, Analog Comparator Positive Input. When configured as an input (DDB0 is cleared [zero]) and with the internal MOS pull-up resistor switched off (PB0 is cleared [zero]), this pin also serves as the positive input of the On-chip Analog Comparator.

Port B Schematics

Note that all port pins are synchronized. The synchronization latches are, however, not shown in the figures.

Figure 22. Port B Schematic Diagram (Pins PB0 and PB1)

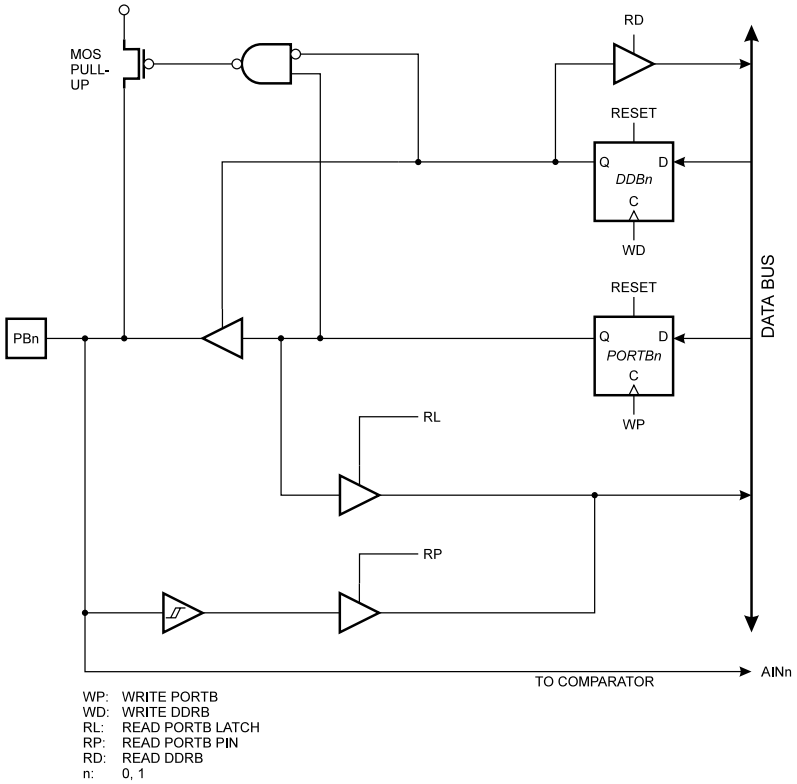


Figure 23. Port B Schematic Diagram (Pins PB2, PB3, and PB4)

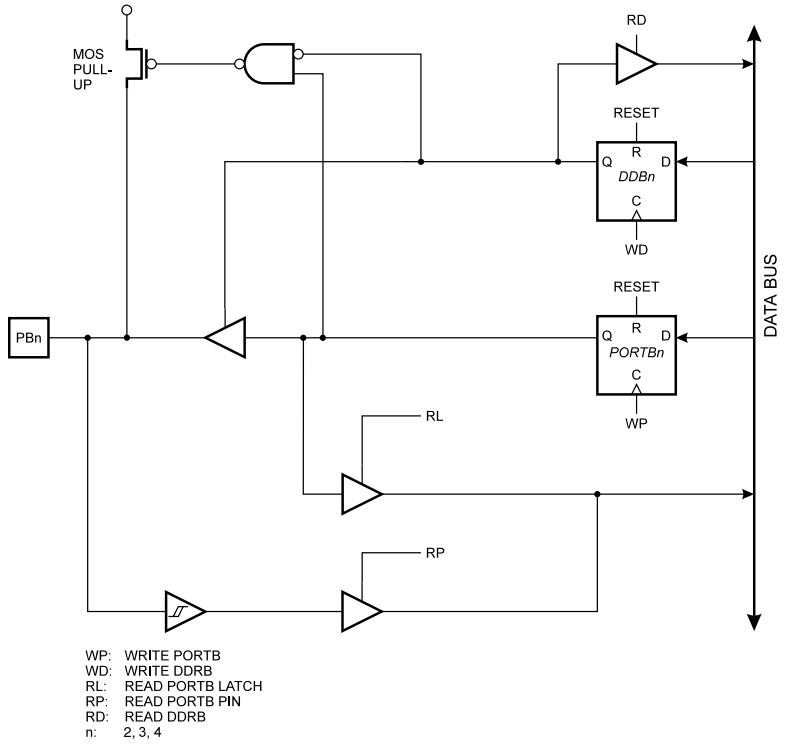


Figure 24. Port B Schematic Diagram (Pin PB5)

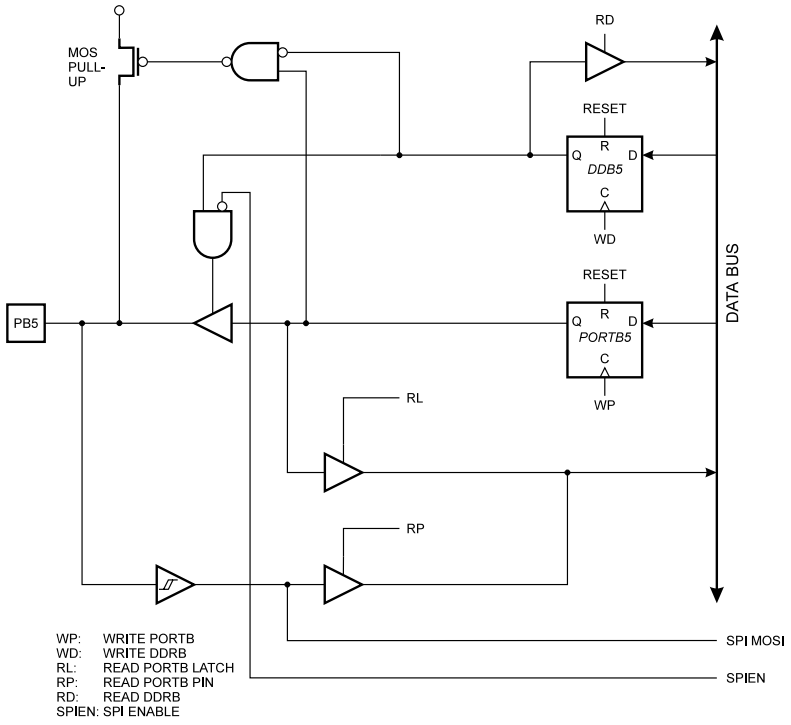


Figure 25. Port B Schematic Diagram (Pin PB6)

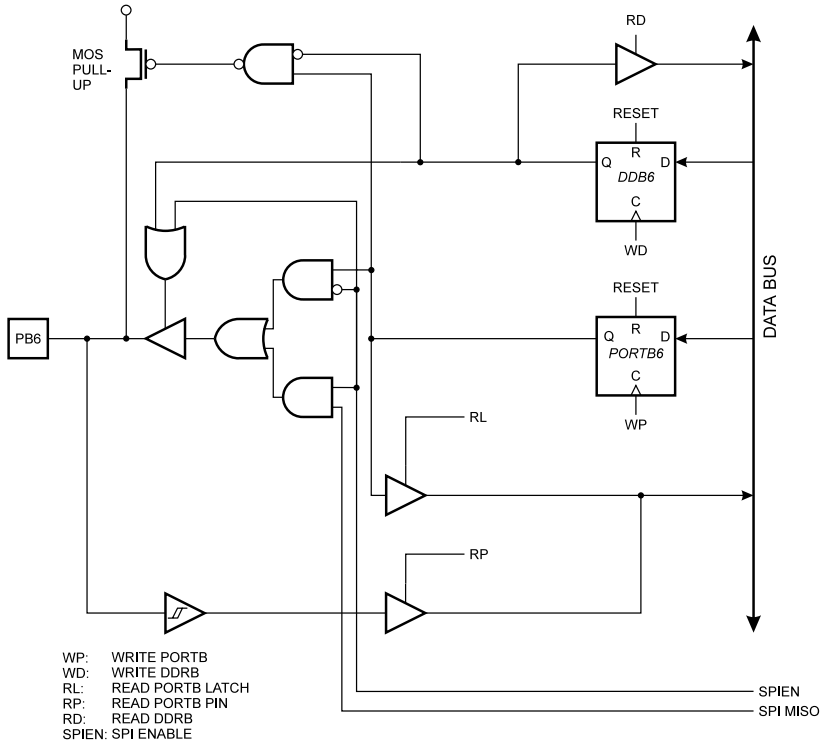
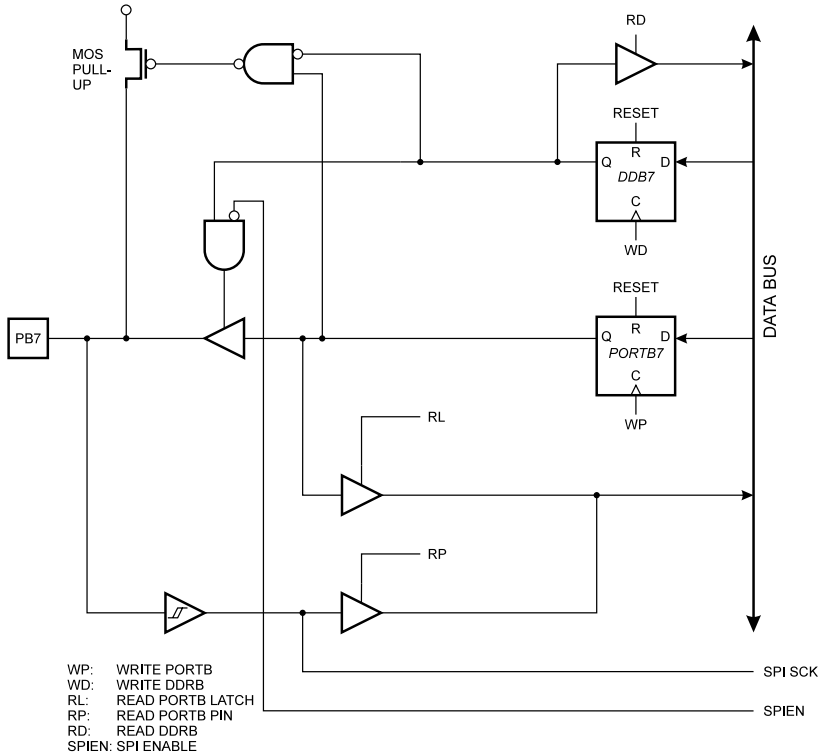


Figure 26. Port B Schematic Diagram (Pin PB7)



Port D

Three I/O memory address locations are allocated for Port D, one each for the Data Register – PORTD ($12), Data Direction Register – DDRD ($11), and the Port D Input Pins – PIND ($10). The Port D Input Pins address is read-only, while the Data Register and the Data Direction Register are read/write.

Port D has seven bi-directional I/O pins with internal pull-up resistors, PD6..PD0. The Port D output buffers can sink 20 mA. As inputs, Port D pins that are externally pulled low will source current if the pull-up resistors are activated.

Some Port D pins have alternate functions as shown in Table 10.

Table 10. Port D Pin Alternate Functions

Port Pin	Alternate Function
PD2	INT0 (External Interrupt 0 input)
PD4	T0 (Timer/Counter 0 external input)

Port D Data Register – PORTD

Bit	7	6	5	4	3	2	1	0	
$12	–	PORTD6	PORTD5	PORTD4	PORTD3	PORTD2	PORTD1	PORTD0	PORTD
Read/Write	R	R/W	R/W	R/W	R/W	R/W	R/W	R/W	
Initial Value	0	0	0	0	0	0	0	0	

Port D Data Direction Register – DDRD

Bit	7	6	5	4	3	2	1	0	
$11	–	DDD6	DDD5	DDD4	DDD3	DDD2	DDD1	DDD0	DDRD
Read/Write	R	R/W	R/W	R/W	R/W	R/W	R/W	R/W	
Initial Value	0	0	0	0	0	0	0	0	

Port D Input Pins Address – PIND

Bit	7	6	5	4	3	2	1	0	
$10	–	PIND6	PIND5	PIND4	PIND3	PIND2	PIND1	PIND0	PIND
Read/Write	R	R	R	R	R	R	R	R	
Initial Value	0	N/A	N/A	N/A	N/A	N/A	N/A	N/A	

The Port D Input Pins address (PIND) is not a register, and this address enables access to the physical value on each Port D pin. When reading PORTD, the Port D Data Latch is read; and when reading PIND, the logical values present on the pins are read.

Port D as General Digital I/O

PDn, general I/O pin: The DDDn bit in the DDRD Register selects the direction of this pin. If DDDn is set (one), PDn is configured as an output pin. If DDDn is cleared (zero), PDn is configured as an input pin. If PORTDn is set (one) when DDDn is configured as an input pin, the MOS pull-up resistor is activated. To switch the pull-up resistor off, the PORTDn bit has to be cleared (zero) or the pin has to be configured as an output pin. The Port D pins are tri-stated when a reset condition becomes active, even if the clock is not active.

Table 11. DDDn Bits' Effect on Port D Pins

DDDn	PORTDn	I/O	Pull-up	Comment
0	0	Input	No	Tri-state (High-Z)
0	1	Input	Yes	PDn will source current if ext. pulled low.
1	0	Output	No	Push-pull Zero Output
1	1	Output	No	Push-pull One Output

Note: n: 6…0, pin number.

Alternate Functions for Port D

The alternate functions of Port D are:

- **T0 – Port D, Bit 4**

T0, Timer/Counter0 clock source. See the timer description for further details.

- **INT0 – Port D, Bit 2**

INT0, External Interrupt source 0. See the interrupt description for further details.

Port D Schematics

Note that all port pins are synchronized. The synchronization latches are, however, not shown in the figures.

Figure 27. Port D Schematic Diagram (Pins PD0, PD1, PD3, PD5, and PD6)

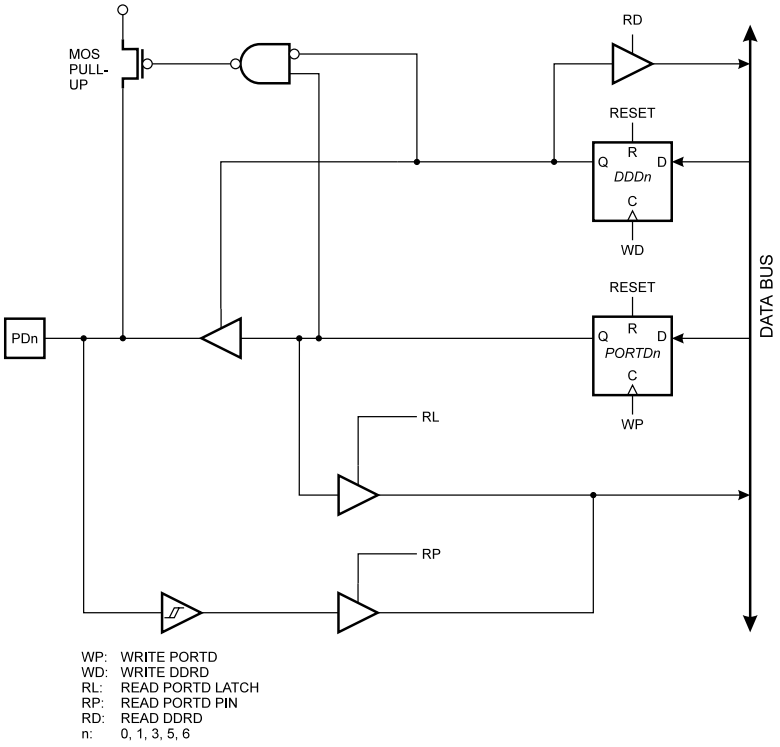


Figure 28. Port D Schematic Diagram (Pin PD2)

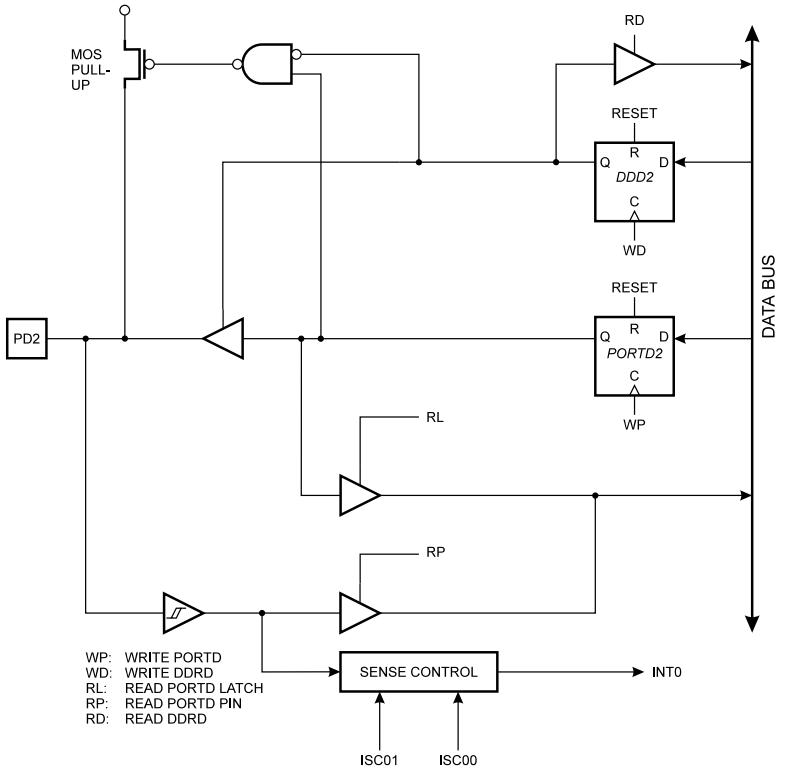
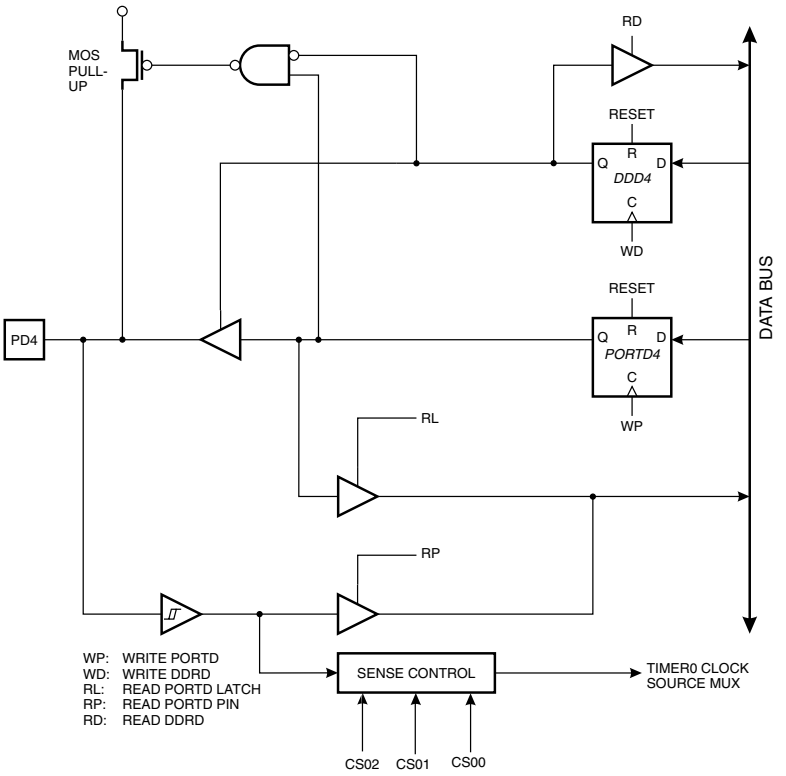


Figure 29. Port D Schematic Diagram (Pin PD4)



Memory Programming

Program and Data Memory Lock Bits

The AT90S1200 MCU provides two Lock bits that can be left unprogrammed ("1") or can be programmed ("0") to obtain the additional features listed in Table 12. The Lock bits can only be erased with the Chip Erase command.

Table 12. Lock Bit Protection Modes

Memory Lock Bits			
Mode	**LB1**	**LB2**	**Protection Type**
1	1	1	No memory lock features enabled.
2	0	1	Further programming of the Flash and EEPROM is disabled.(1)
3	0	0	Same as mode 2, and verify is also disabled.

Note: 1. In Parallel mode, further programming of the Fuse bits are also disabled. Program the Fuse bits before programming the Lock bits.

Fuse Bits

The AT90S1200 has two Fuse bits: SPIEN and RCEN.

- When the SPIEN Fuse bit is programmed ("0"), Serial Program Downloading is enabled. Default value is programmed ("0").
- When the RCEN Fuse bit is programmed ("0"), MCU clocking from the Internal RC Oscillator is selected. Default value is erased ("1"). Parts with this bit pre-programmed ("0") can be delivered on demand.
- The Fuse bits are not accessible in Serial Programming mode. The status of the Fuse bits is not affected by Chip Erase.

Signature Bytes

All Atmel microcontrollers have a 3-byte signature code that identifies the device. This code can be read in both Serial and Parallel modes. The three bytes reside in a separate address space.

For the AT90S1200 they are:

1. $00: $1E (indicates manufactured by Atmel)
2. $01: $90 (indicates 1 Kb Flash memory)
3. $02: $01 (indicates AT90S1200 device when $01 is $90)

Note: When both Lock bits are programmed (lock mode 3), the signature bytes cannot be read in Serial mode. Reading the signature bytes will return: $00, $01 and $02.

Programming the Flash and EEPROM

Atmel's AT90S1200 offers 1K byte of in-System Reprogrammable Flash program memory and 64 bytes of EEPROM data memory.

The AT90S1200 is normally shipped with the On-chip Flash program memory and EEPROM data memory arrays in the erased state (i.e., contents = $FF) and ready to be programmed. This device supports a High-voltage (12V) Parallel Programming mode and a Low-voltage Serial Programming mode. The +12V is used for programming enable only, and no current of significance is drawn by this pin. The Serial Programming mode provides a convenient way to download program and data into the AT90S1200 inside the user's system.

The program and data memory arrays on the AT90S1200 are programmed byte-by-byte in either programming mode. For the EEPROM, an auto-erase cycle is provided within

the self-timed write instruction in the Serial Programming mode. During programming, the supply voltage must be in accordance with Table 13.

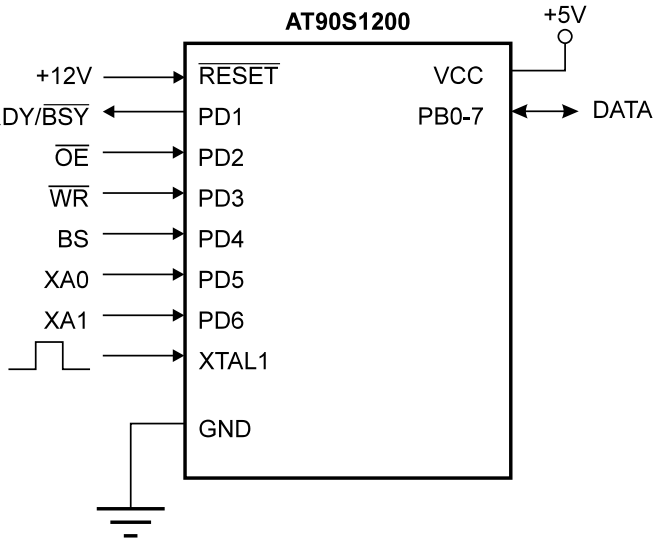
Table 13. Supply Voltage during Programming

Part	Serial Programming	Parallel Programming
AT90S1200	2.7 - 6.0V	4.5 - 5.5V

Parallel Programming

This section describes how to parallel program and verify Flash program memory, EEPROM data memory, Lock bits and Fuse bits in the AT90S1200.

Figure 30. Parallel Programming



Signal Names

In this section, some pins of the AT90S1200 are referenced by signal names describing their function during parallel programming rather than their pin names, see Figure 30 and Table 14. Pins not described in Table 14 are referenced by pin names.

The XA1/XA0 pins determines the action executed when the XTAL1 pin is given a positive pulse. The coding is shown in Table 15.

When pulsing $\overline{WR}$ or $\overline{OE}$, the command loaded determines the action executed. The command is a byte where the different bits are assigned functions as shown in Table 16.

Table 14. Pin Name Mapping

Signal Name in Programming Mode	Pin Name	I/O	Function
RDY/$\overline{BSY}$	PD1	O	0: Device is busy programming, 1: Device is ready for new command
$\overline{OE}$	PD2	I	Output Enable (Active low)
$\overline{WR}$	PD3	I	Write Pulse (Active low)
BS	PD4	I	Byte Select (“0” selects low byte, “1” selects high byte)
XA0	PD5	I	XTAL Action Bit 0
XA1	PD6	I	XTAL Action Bit 1
DATA	PB0-7	I/O	Bi-directional Data Bus (Output when $\overline{OE}$ is low)

Table 15. XA1 and XA0 Coding

XA1	XA0	Action when XTAL1 is Pulsed
0	0	Load Flash or EEPROM Address (High or low address byte for Flash determined by BS).
0	1	Load Data (High or low data byte for Flash determined by BS).
1	0	Load Command
1	1	No Action, Idle

Table 16. Command Byte Coding

Command Byte	Command Executed
1000 0000	Chip Erase
0100 0000	Write Fuse Bits
0010 0000	Write Lock Bits
0001 0000	Write Flash
0001 0001	Write EEPROM
0000 1000	Read Signature Bytes
0000 0100	Read Fuse and Lock Bits
0000 0010	Read Flash
0000 0011	Read EEPROM

Enter Programming Mode

The following algorithm puts the device in Parallel Programming mode:

1. Apply supply voltage according to Table 13, between V_{CC} and GND.
2. Set the $\overline{\text{RESET}}$ and BS pin to “0” and wait at least 100 ns.
3. Apply 11.5 - 12.5V to $\overline{\text{RESET}}$. Any activity on BS within 100 ns after +12V has been applied to $\overline{\text{RESET}}$, will cause the device to fail entering Programming mode.

Chip Erase

The Chip Erase command will erase the Flash and EEPROM memories, and the Lock bits. The Lock bits are not Reset until the Flash and EEPROM have been completely erased. The Fuse bits are not changed. Chip Erase must be performed before the Flash or EEPROM is reprogrammed.

Load Command “Chip Erase”

1. Set XA1, XA0 to “10”. This enables command loading.
2. Set BS to “0”.
3. Set DATA to “1000 0000”. This is the command for Chip Erase.
4. Give XTAL1 a positive pulse. This loads the command.
5. Give $\overline{\text{WR}}$ a t_{WLWH_CE} wide negative pulse to execute Chip Erase, t_{WLWH_CE} is found in Table 17. Chip Erase does not generate any activity on the RDY/$\overline{\text{BSY}}$ pin.

Programming the Flash

A: Load Command “Write Flash”

1. Set XA1, XA0 to “10”. This enables command loading.
2. Set BS to “0”.
3. Set DATA to “0001 0000”. This is the command for Write Flash.

4. Give XTAL1 a positive pulse. This loads the command.

B: Load Address High Byte

1. Set XA1, XA0 to “00”. This enables address loading.
2. Set BS to “1”. This selects high byte.
3. Set DATA = Address high byte ($00 - $01).
4. Give XTAL1 a positive pulse. This loads the address high byte.

C: Load Address Low Byte

1. Set XA1, XA0 to “00”. This enables address loading.
2. Set BS to “0”. This selects low byte.
3. Set DATA = Address low byte ($00 - $FF).
4. Give XTAL1 a positive pulse. This loads the address low byte.

D: Load Data Low Byte

1. Set XA1, XA0 to “01”. This enables data loading.
2. Set DATA = Data low byte ($00 - $FF).
3. Give XTAL1 a positive pulse. This loads the data low byte.

E: Write Data Low Byte

1. Set BS to “0”. This selects low data.
2. Give $\overline{WR}$ a negative pulse. This starts programming of the data byte. RDY/$\overline{BSY}$ goes low.
3. Wait until RDY/$\overline{BSY}$ goes high to program the next byte.

(See Figure 31 for signal waveforms.)

F: Load Data High Byte

1. Set XA1, XA0 to “01”. This enables data loading.
2. Set DATA = Data high byte ($00 - $FF).
3. Give XTAL1 a positive pulse. This loads the data high byte.

G: Write Data High Byte

1. Set BS to “1”. This selects high data.
2. Give $\overline{WR}$ a negative pulse. This starts programming of the data byte. RDY/$\overline{BSY}$ goes low.
3. Wait until RDY/$\overline{BSY}$ goes high to program the next byte.

(See Figure 32 for signal waveforms.)

The loaded command and address are retained in the device during programming. For efficient programming, the following should be considered:

- The command needs only be loaded once when writing or reading multiple memory locations.
- Address high byte needs only be loaded before programming a new 256-word page in the Flash.
- Skip writing the data value $FF; that is, the contents of the entire Flash and EEPROM after a Chip Erase.

These considerations also apply to EEPROM programming and Flash, EEPROM and signature byte reading.

Figure 31. Programming the Flash Waveforms

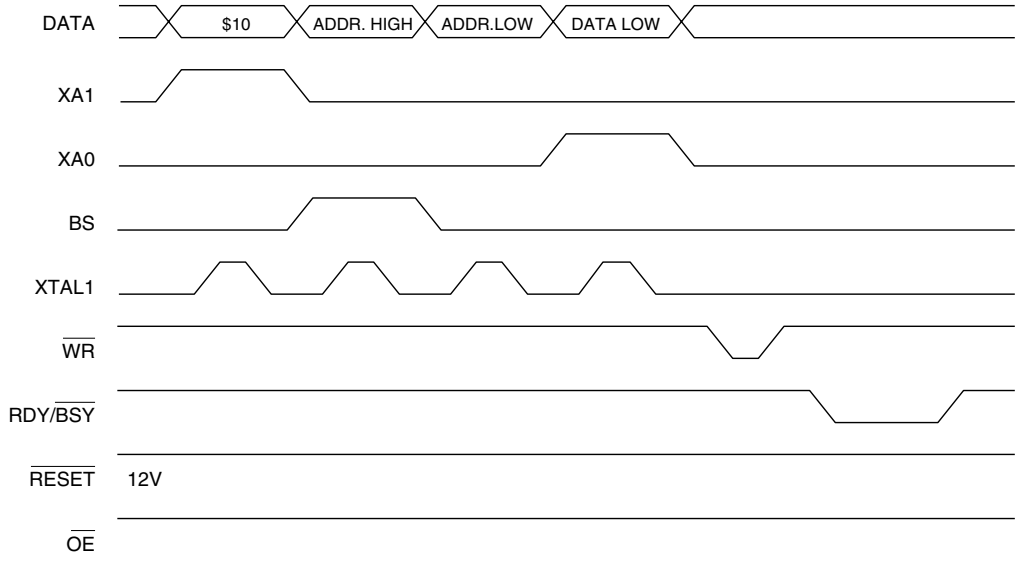
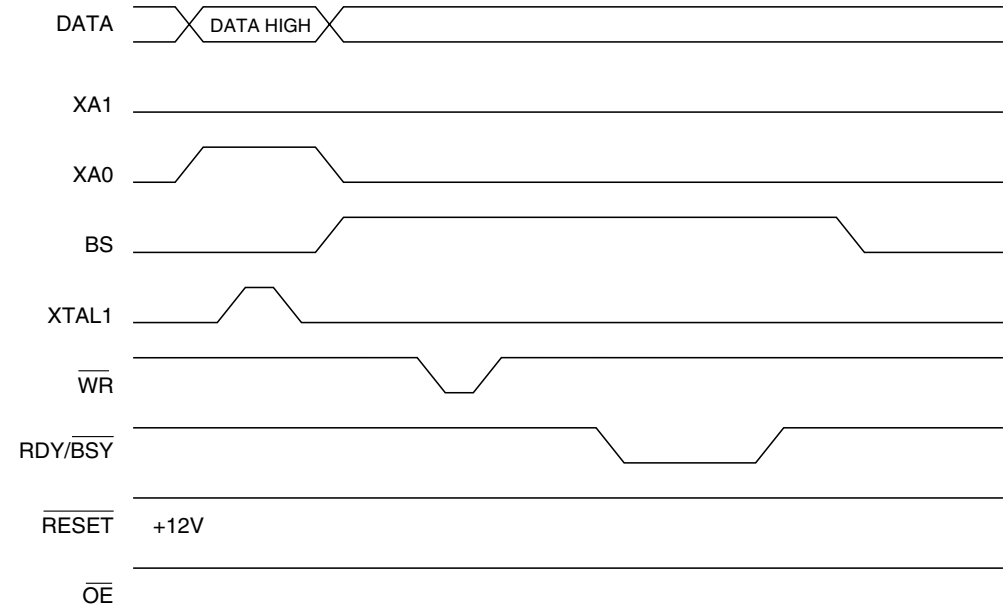


Figure 32. Programming the Flash Waveforms (Continued)



Reading the Flash

The algorithm for reading the Flash memory is as follows (refer to “Programming the Flash” for details on command and address loading):

1. A: Load Command “0000 0010”.
2. B: Load Address High Byte ($00 - $01).
3. C: Load Address Low Byte ($00 - $FF).
4. Set $\overline{OE}$ to “0”, and BS to “0”. The Flash word low byte can now be read at DATA.
5. Set BS to “1”. The Flash word high byte can now be read from DATA.
6. Set $\overline{OE}$ to “1”.

Programming the EEPROM

The programming algorithm for the EEPROM data memory is as follows (refer to “Programming the Flash” for details on command, address and data loading):

1. A: Load Command “0001 0001”.
2. C: Load Address Low Byte ($00 - $3F).
3. D: Load Data Low Byte ($00 - $FF).
4. E: Write Data Low Byte.

Reading the EEPROM

The algorithm for reading the EEPROM memory is as follows (refer to “Programming the Flash” for details on command and address loading):

1. A: Load Command “0000 0011”.
2. C: Load Address Low Byte ($00 - $3F).
3. Set $\overline{OE}$ to “0”, and BS to “0”. The EEPROM data byte can now be read at DATA.
4. Set $\overline{OE}$ to “1”.

Programming the Fuse Bits

The algorithm for programming the Fuse bits is as follows (refer to “Programming the Flash” for details on command and data loading):

1. A: Load Command “0100 0000”.
2. D: Load Data Low Byte. Bit n = “0” programs and bit n = “1” erases the Fuse bit.

   Bit 5 = SPIEN Fuse
   Bit 0 = RCEN Fuse

   Bit 7 - 6, 4 - 1 = “1”. These bits are reserved and should be left unprogrammed (“1”).
3. Give $\overline{WR}$ a t_{WLWH_PFB} wide negative pulse to execute the programming; t_{WLWH_PFB} is found in Table 17. Programming the Fuse bits does not generate any activity on the RDY/$\overline{BSY}$ pin.

Programming the Lock Bits

The algorithm for programming the Lock bits is as follows (refer to “Programming the Flash” for details on command and data loading):

1. A: Load Command “0010 0000”.
2. D: Load Data Low Byte. Bit n = “0” programs the Lock bit.

   Bit 2 = Lock Bit2
   Bit 1 = Lock Bit1

   Bit 7 - 3, 0 = “1”. These bits are reserved and should be left unprogrammed (“1”).
3. E: Write Data Low Byte.

The Lock bits can only be cleared by executing Chip Erase.

Reading the Fuse and Lock Bits

The algorithm for reading the Fuse and Lock bits is as follows (refer to “Programming the Flash” on page 39 for details on command loading):

1. A: Load Command “0000 0100”.
2. Set $\overline{OE}$ to “0”, and BS to “1”. The status of Fuse and Lock bits can now be read at DATA (“0” means programmed).

   Bit 7 = Lock Bit1
   Bit 6 = Lock Bit2
   Bit 5 = SPIEN Fuse
   Bit 0 = RCEN Fuse
3. Set $\overline{OE}$ to “1”.

Observe especially that BS needs to be set to “1”.

Reading the Signature Bytes

The algorithm for reading the signature bytes is as follows (refer to “Programming the Flash” on page 39 for details on command and address loading):

1. A: Load Command “0000 1000”.
2. C: Load Address Low Byte ($00 - $02).
   Set $\overline{OE}$ to “0”, and BS to “0”. The selected signature byte can now be read at DATA.
   Set $\overline{OE}$ to “1”.

Parallel Programming Characteristics

Figure 33. Parallel Programming Timing

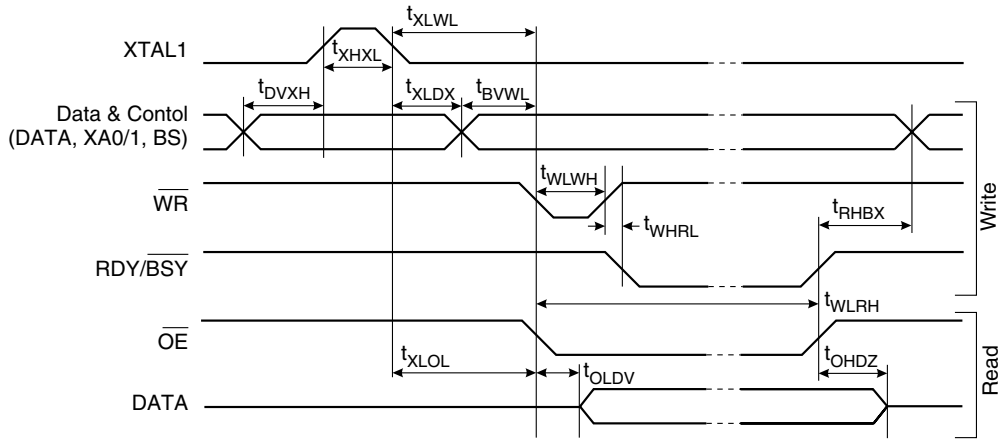


Table 17. Parallel Programming Characteristics, T_A = 25°C ± 10%, V_{CC} = 5V ± 10%

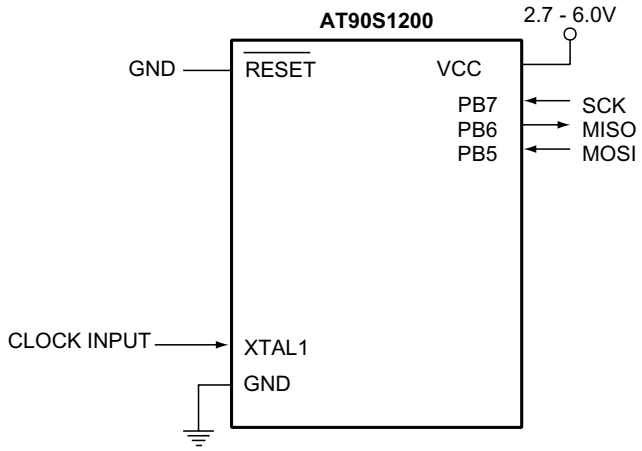
Symbol	Parameter	Min	Typ	Max	Units
V_{PP}	Programming Enable Voltage	11.5		12.5	V
I_{PP}	Programming Enable Current			250.0	µA
t_{DVXH}	Data and Control Setup before XTAL1 High	67.0			ns
t_{XHXL}	XTAL1 Pulse Width High	67.0			ns
t_{XLDX}	Data and Control Hold after XTAL1 Low	67.0			ns
t_{XLWL}	XTAL1 Low to $\overline{WR}$ Low	67.0			ns
t_{BVWL}	BS Valid to $\overline{WR}$ Low	67.0			ns
t_{RHBX}	BS Hold after RDY/$\overline{BSY}$ High	67.0			ns
t_{WLWH}	$\overline{WR}$ Pulse Width Low(1)	67.0			ns
t_{WHRL}	$\overline{WR}$ High to RDY/$\overline{BSY}$ Low(2)		20.0		ns
t_{WLRH}	$\overline{WR}$ Low to RDY/$\overline{BSY}$ High(2)	0.5	0.7	0.9	ms
t_{XLOL}	XTAL1 Low to $\overline{OE}$ Low	67.0			ns
t_{OLDV}	$\overline{OE}$ Low to DATA Valid		20.0		ns
t_{OHDZ}	$\overline{OE}$ High to DATA Tri-stated			20.0	ns
t_{WLWH_CE}	$\overline{WR}$ Pulse Width Low for Chip Erase	5.0	10.0	15.0	ms
t_{WLWH_PFB}	$\overline{WR}$ Pulse Width Low for Programming the Fuse Bits	1.0	1.5	1.8	ms

Notes: 1. Use t_{WLWH_CE} for chip erase and t_{WLWH_PFB} for programming the Fuse bits.
2. If t_{WLWH} is held longer than t_{WLRH}, no RDY/$\overline{BSY}$ pulse will be seen.

Serial Downloading

Both the program and data memory arrays can be programmed using the SPI bus while $\overline{\text{RESET}}$ is pulled to GND. The serial interface consists of pins SCK, MOSI (input) and MISO (output) (see Figure 34). After $\overline{\text{RESET}}$ is set low, the Programming Enable instruction needs to be executed first before program/erase instructions can be executed.

Figure 34. Serial Programming and Verify



Note: If the device is clocked by the Internal Oscillator, it is no need to connect a clock source to the XTAL1 pin

For the EEPROM, an auto-erase cycle is provided within the self-timed write instruction and there is no need to first execute the Chip Erase instruction. The Chip Erase instruction turns the content of every memory location in both the Program and EEPROM arrays into $FF.

The program and EEPROM memory arrays have separate address spaces: $0000 to $01FF for Flash program memory and $000 to $03F for EEPROM data memory.

Either an external system clock is supplied at pin XTAL1 or a crystal needs to be connected across pins XTAL1 and XTAL2. The minimum low and high periods for the Serial Clock (SCK) input are defined as follows:

Low: > 1 XTAL1 clock cycle

High: > 4 XTAL1 clock cycles

Serial Programming Algorithm

When writing serial data to the AT90S1200, data is clocked on the rising edge of SCK.

When reading data from the AT90S1200, data is clocked on the falling edge of SCK. See Figure 35 and Table 20 for timing details.

To program and verify the AT90S1200 in the Serial Programming mode, the following sequence is recommended (See 4-byte instruction formats in Table 17):

1. Power-up sequence:

   Apply power between V_{CC} and GND while $\overline{\text{RESET}}$ and SCK are set to “0”. If a crystal is not connected across pins XTAL1 and XTAL2 or the device is not running from the Internal RC Oscillator, apply a clock signal to the XTAL1 pin. If the programmer can not guarantee that SCK is held low during power-up, $\overline{\text{RESET}}$ must be given a positive pulse after SCK has been set to “0”.
2. Wait for at least 20 ms and enable serial programming by sending the Programming Enable serial instruction to the MOSI (PB5) pin.

3. If a Chip Erase is performed (must be done to erase the Flash), wait t_{WD_ERASE} after the instruction, give $\overline{\text{RESET}}$ a positive pulse, and start over from step 2. See Table 21 on page 47 for t_{WD_ERASE} value.
4. The Flash or EEPROM array is programmed one byte at a time by supplying the address and data together with the appropriate Write instruction. An EEPROM memory location is first automatically erased before new data is written. Wait t_{WD_PROG} after transmitting the instruction. In an erased device, no $FFs in the data file(s) needs to be programmed. See Table 22 on page 47 for t_{WD_PROG} value.
5. Any memory location can be verified by using the Read instruction which returns the content at the selected address at the serial output MISO (PB6) pin.

   At the end of the programming session, $\overline{\text{RESET}}$ can be set high to commence normal operation.
6. Power-off sequence (if needed):

   Set XTAL1 to “0” (if a crystal is not used or the device is running from the Internal RC Oscillator).

   Set $\overline{\text{RESET}}$ to “1”.

   Turn V_{CC} power off.

Data Polling EEPROM

When a byte is being programmed into the EEPROM, reading the address location being programmed will give the value P1 until the auto-erase is finished, and then the value P2. See Table 18 for P1 and P2 values.

At the time the device is ready for a new EEPROM byte, the programmed value will read correctly. This is used to determine when the next byte can be written. This will not work for the values P1 and P2, so when programming these values, the user will have to wait for at least the prescribed time t_{WD_PROG} before programming the next byte. See Table 22 for t_{WD_PROG} value. As a chip-erased device contains $FF in all locations, programming of addresses that are meant to contain $FF can be skipped. This does not apply if the EEPROM is reprogrammed without first chip-erasing the device.

Table 18. Read Back Value during EEPROM Polling

Part	P1	P2
AT90S1200	$00	$FF

Data Polling Flash

When a byte is being programmed into the Flash, reading the address location being programmed will give the value $FF. At the time the device is ready for a new byte, the programmed value will read correctly. This is used to determine when the next byte can be written. This will not work for the value $FF, so when programming this value, the user will have to wait for at least t_{WD_PROG} before programming the next byte. As a chip-erased device contains $FF in all locations, programming of addresses that are meant to contain $FF, can be skipped.

Figure 35. Serial Programming Waveforms

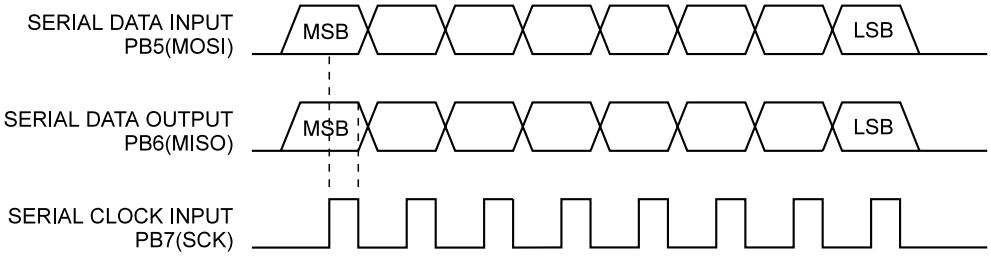


Table 19. Serial Programming Instruction Set for AT90S1200

Instruction	Instruction Format				Operation
	Byte 1	Byte 2	Byte 3	Byte4	
Programming Enable	1010 1100	0101 0011	xxxx xxxx	xxxx xxxx	Enable serial programming while $\overline{\text{RESET}}$ is low.
Chip Erase	1010 1100	100x xxxx	xxxx xxxx	xxxx xxxx	Chip erase both Flash and EEPROM memory arrays.
Read Program Memory	0010 **H**000	0000 000**a**	**bbbb bbbb**	**oooo oooo**	Read **H** (high or low) byte **o** from program memory at word address **a**:**b**.
Write Program Memory	0100 **H**000	0000 000**a**	**bbbb bbbb**	**iiii iiii**	Write **H** (high or low) byte **i** to program memory at word address **a**:**b**.
Read EEPROM Memory	1010 0000	0000 0000	00**bb bbbb**	**oooo oooo**	Read data **o** from EEPROM memory at address **b**.
Write EEPROM Memory	1100 0000	0000 0000	00**bb bbbb**	**iiii iiii**	Write data **i** to EEPROM memory at address **b**.
Write Lock Bits	1010 1100	1111 1***21***1	xxxx xxxx	xxxx xxxx	Write Lock bits. Set bits ***1***,***2*** = “0” to program Lock bits.
Read Signature Byte	0011 0000	xxxx xxxx	xxxx xx**bb**	**oooo oooo**	Read signature byte **o** from address **b**.(1)

Note: **a** = address high bits, **b** = address low bits, **H** = 0 – Low byte, 1 – High byte, **o** = data out, **i** = data in, x = don’t care, ***1*** = Lock Bit 1, ***2*** = Lock Bit 2

Note: 1. The signature bytes are not readable in lock mode 3 (i.e., both Lock bits programmed).

Serial Programming Characteristics

Figure 36. Serial Programming Timing

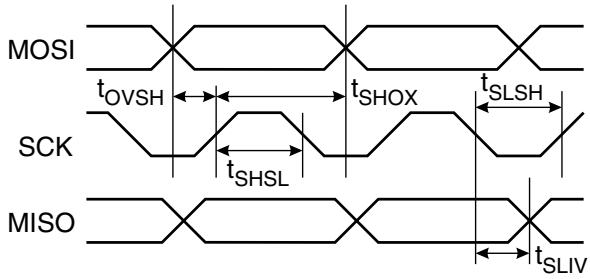


Table 20. Serial Programming Characteristics, T_A = -40°C to 85°C, V_{CC} = 2.7 - 6.0V (unless otherwise noted)

Symbol	Parameter	Min	Typ	Max	Units
$1/t_{CLCL}$	Oscillator Frequency (V_{CC} = 2.7 - 4.0V)	0		4.0	MHz
t_{CLCL}	Oscillator Period (V_{CC} = 2.7 - 4.0V)	250.0			ns
$1/t_{CLCL}$	Oscillator Frequency (V_{CC} = 4.0 - 6.0V)	0		12.0	MHz
t_{CLCL}	Oscillator Period (V_{CC} = 4.0 - 6.0V)	83.3			ns
t_{SHSL}	SCK Pulse Width High	4.0 t_{CLCL}			ns
t_{SLSH}	SCK Pulse Width Low	t_{CLCL}			ns
t_{OVSH}	MOSI Setup to SCK High	1.25 t_{CLCL}			ns
t_{SHOX}	MOSI Hold after SCK High	2.5 t_{CLCL}			ns
t_{SLIV}	SCK Low to MISO Valid	10.0	16.0	32.0	ns

Table 21. Minimum Wait Delay after the Chip Erase Instruction

Symbol	3.2V	3.6V	4.0V	5.0V
t_{WD_ERASE}	18 ms	14 ms	12 ms	8 ms

Table 22. Minimum Wait Delay after Writing a Flash or EEPROM Location

Symbol	3.2V	3.6V	4.0V	5.0V
t_{WD_PROG}	9 ms	7 ms	6 ms	4 ms

Electrical Characteristics

Absolute Maximum Ratings*

Operating Temperature	-55°C to +125°C
Storage Temperature	-65°C to +150°C
Voltage on Any Pin Except $\overline{\text{RESET}}$ with Respect to Ground	-1.0V to V_{CC}+0.5V
Voltage on $\overline{\text{RESET}}$ with Respect to Ground	-1.0V to +13.0V
Maximum Operating Voltage	6.6V
DC Current per I/O Pin	40.0 mA
DC Current V_{CC} and GND Pins	200.0 mA

*NOTICE: Stresses beyond those listed under "Absolute Maximum Ratings" may cause permanent damage to the device. This is a stress rating only and functional operation of the device at these or other conditions beyond those indicated in the operational sections of this specification is not implied. Exposure to absolute maximum rating conditions for extended periods may affect device reliability.

DC Characteristics

T_A = -40×C to 85×C, V_{CC} = 2.7V to 6.0V (unless otherwise noted)

Symbol	Parameter	Condition	Min	Typ	Max	Units
V_{IL}	Input Low Voltage	(Except XTAL1)	-0.5		0.3 V_{CC}(1)	V
V_{IL1}	Input Low Voltage	(XTAL1)	-0.5		0.3 V_{CC}(1)	V
V_{IH}	Input High Voltage	(Except XTAL1, $\overline{\text{RESET}}$)	0.6 V_{CC}(2)		V_{CC} + 0.5	V
V_{IH1}	Input High Voltage	(XTAL1)	0.7 V_{CC}(2)		V_{CC} + 0.5	V
V_{IH2}	Input High Voltage	($\overline{\text{RESET}}$)	0.85 V_{CC}(2)		V_{CC} + 0.5	V
V_{OL}	Output Low Voltage(3) (Ports B, D)	I_{OL} = 20 mA, V_{CC} = 5V I_{OL} = 10 mA, V_{CC} = 3V			0.6 0.5	V V
V_{OH}	Output High Voltage(4) (Ports B, D)	I_{OH} = -3 mA, V_{CC} = 5V I_{OH} = -1.5 mA, V_{CC} = 3V	4.3 2.3			V V
I_{IL}	Input Leakage Current I/O pin	V_{CC} = 6V, pin low (absolute value)			8.0	µA
I_{IH}	Input Leakage Current I/O pin	V_{CC} = 6V, pin high (absolute value)			980.0	nA
RRST	Reset Pull-up Resistor		100.0		500.0	kΩ
$R_{I/O}$	I/O Pin Pull-up Resistor		35.0		120.0	kΩ
I_{CC}	Power Supply Current	Active Mode, V_{CC} = 3V, 4 MHz			3.0	mA
		Idle Mode V_{CC} = 3V, 4 MHz			1.0	mA
I_{CC}	Power-down mode(5)	WDT enabled, V_{CC} = 3V		9.0	15.0	µA
		WDT disabled, V_{CC} = 3V		<1.0	2.0	µA

DC Characteristics

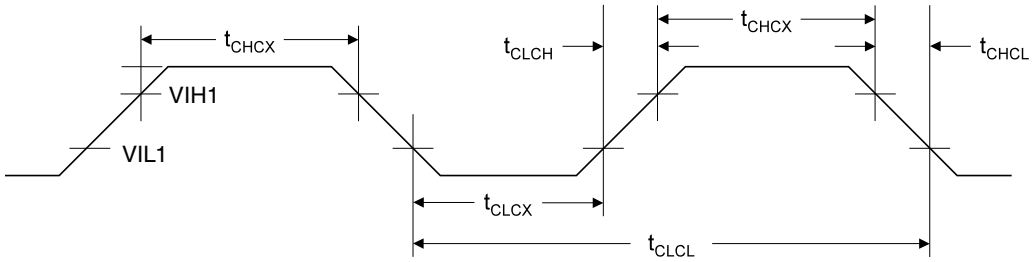
T_A = -40×C to 85×C, V_{CC} = 2.7V to 6.0V (unless otherwise noted) (Continued)

Symbol	Parameter	Condition	Min	Typ	Max	Units
V_{ACIO}	Analog Comparator Input Offset Voltage	V_{CC} = 5V V_{in} = V_{CC}/ 2			40.0	mV
I_{ACLK}	Analog Comparator Input Leakage Current	V_{CC} = 5V V_{in} = V_{CC}/ 2	-50.0		50.0	nA
t_{ACPD}	Analog Comparator Propagation Delay	V_{CC} = 2.7V V_{CC} = 4.0V		750.0 500.0		ns

Notes:
1. “Max” means the highest value where the pin is guaranteed to be read as low.
2. “Min” means the lowest value where the pin is guaranteed to be read as high.
3. Although each I/O port can sink more than the test conditions (20 mA at V_{CC} = 5V, 10 mA at V_{CC} = 3V) under steady state conditions (non-transient), the following must be observed:
   1] The sum of all I_{OL}, for all ports, should not exceed 200 mA.
   2] The sum of all I_{OL}, for port D0 - D5 and XTAL2, should not exceed 100 mA.
   3] The sum of all I_{OL}, for ports B0 - B7 and D6, should not exceed 100 mA.
   If I_{OL} exceeds the test condition, V_{OL} may exceed the related specification. Pins are not guaranteed to sink current greater than the listed test condition.
4. Although each I/O port can source more than the test conditions (3 mA at V_{CC} = 5V, 1.5 mA at V_{CC} = 3V) under steady state conditions (non-transient), the following must be observed:
   1] The sum of all I_{OH}, for all ports, should not exceed 200 mA.
   2] The sum of all I_{OH}, for port D0 - D5 and XTAL2, should not exceed 100 mA.
   3] The sum of all I_{OH}, for ports B0 - B7 and D6, should not exceed 100 mA.
   If I_{OH} exceeds the test condition, V_{OH} may exceed the related specification. Pins are not guaranteed to source current greater than the listed test condition.
5. Minimum V_{CC} for power-down is 2V.

External Clock Drive Waveforms

Figure 37. External Clock Drive



External Clock Drive

Table 23. External Clock Drive

Symbol	Parameter	V_{CC} = 2.7V to 4.0V		V_{CC} = 4.0V to 6.0V		Units
		Min	Max	Min	Max	
$1/t_{CLCL}$	Oscillator Frequency	0	4.0	0	12.0	MHz
t_{CLCL}	Clock Period	250.0		83.3		ns
t_{CHCX}	High Time	100.0		33.3		ns
t_{CLCX}	Low Time	100.0		33.3		ns
t_{CLCH}	Rise Time		1.6		0.5	µs
t_{CHCL}	Fall Time		1.6		0.5	µs

Typical Characteristics

The following charts show typical behavior. These figures are not tested during manufacturing. All current consumption measurements are performed with all I/O pins configured as inputs and with internal pull-ups enabled. A sine wave generator with rail-to-rail output is used as clock source.

The power consumption in Power-down mode is independent of clock selection.

The current consumption is a function of several factors such as: operating voltage, operating frequency, loading of I/O pins, switching rate of I/O pins, code executed and ambient temperature. The dominating factors are operating voltage and frequency.

The current drawn from capacitive loaded pins may be estimated (for one pin) as $C_L \bullet V_{CC} \bullet f$ where C_L = load capacitance, V_{CC} = operating voltage and f = average switching frequency of I/O pin.

The parts are characterized at frequencies higher than test limits. Parts are not guaranteed to function properly at frequencies higher than the ordering code indicates.

The difference between current consumption in Power-down mode with Watchdog Timer enabled and Power-down mode with Watchdog Timer disabled represents the differential current drawn by the Watchdog Timer.

Figure 38. Active Supply Current vs. Frequency

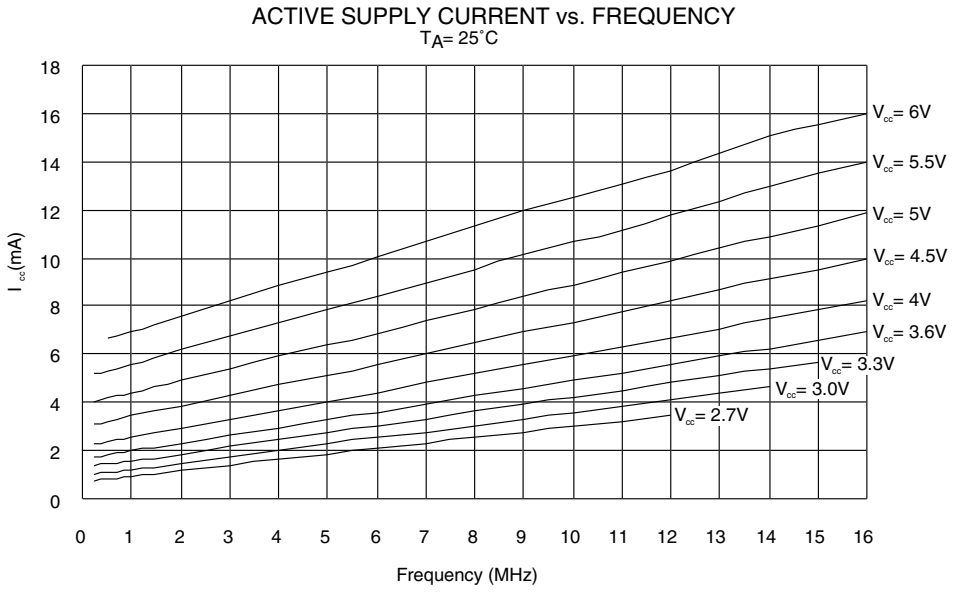


Figure 39. Active Supply Current vs. V_{CC}

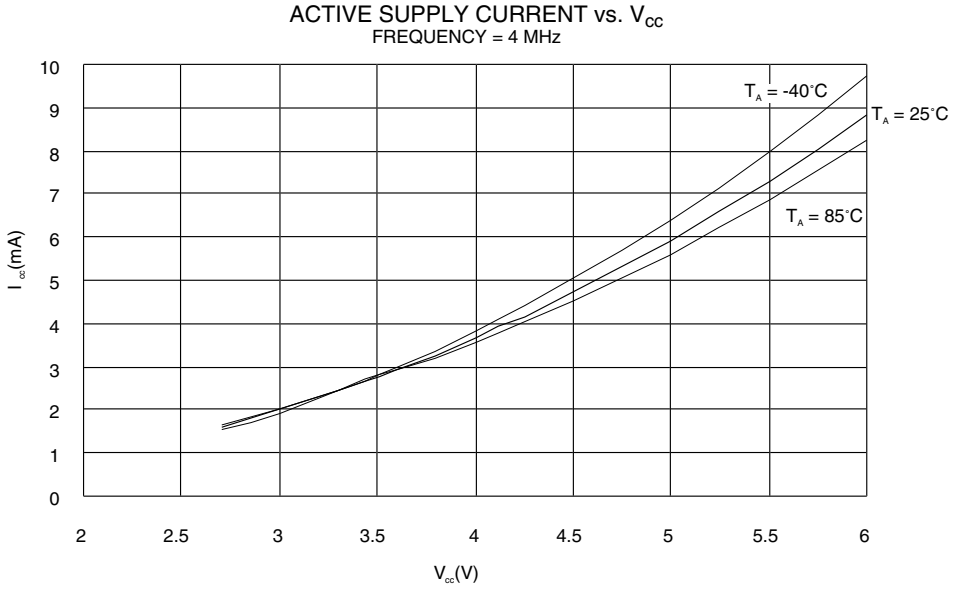


Figure 40. Active Supply Current vs. V_{CC}, Device Clocked by Internal Oscillator

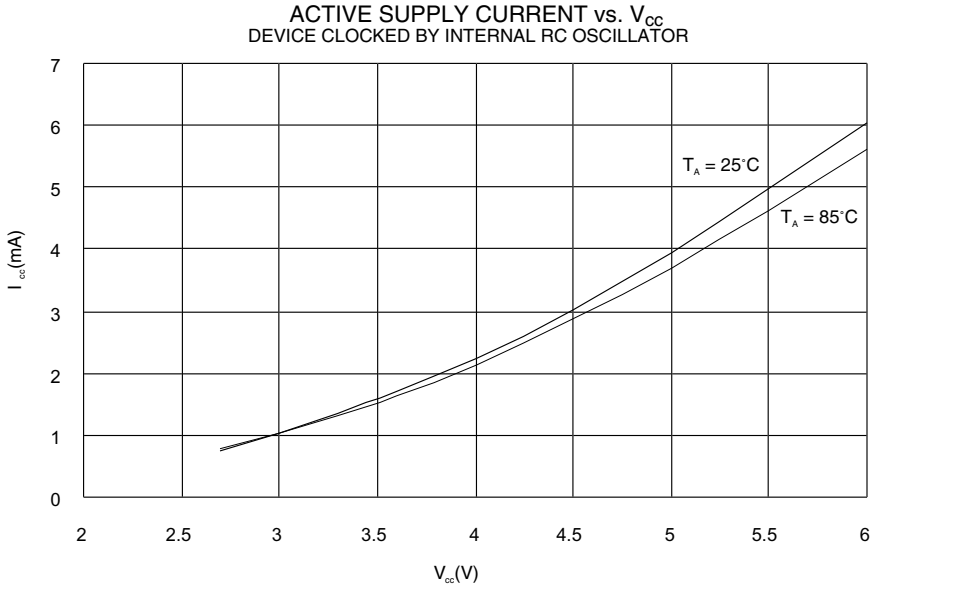


Figure 41. Idle Supply Current vs. Frequency

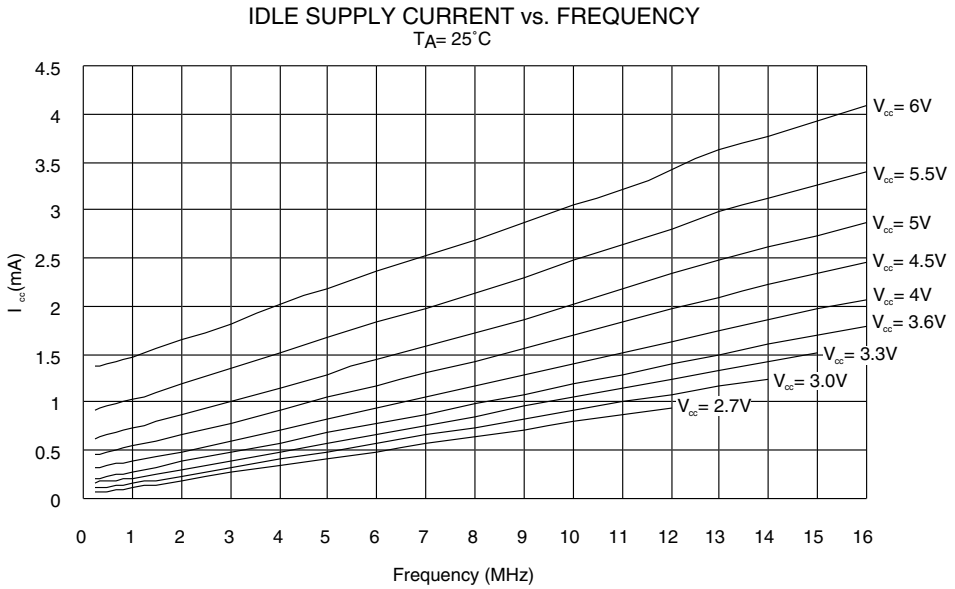


Figure 42. Idle Supply Current vs. V_{CC}

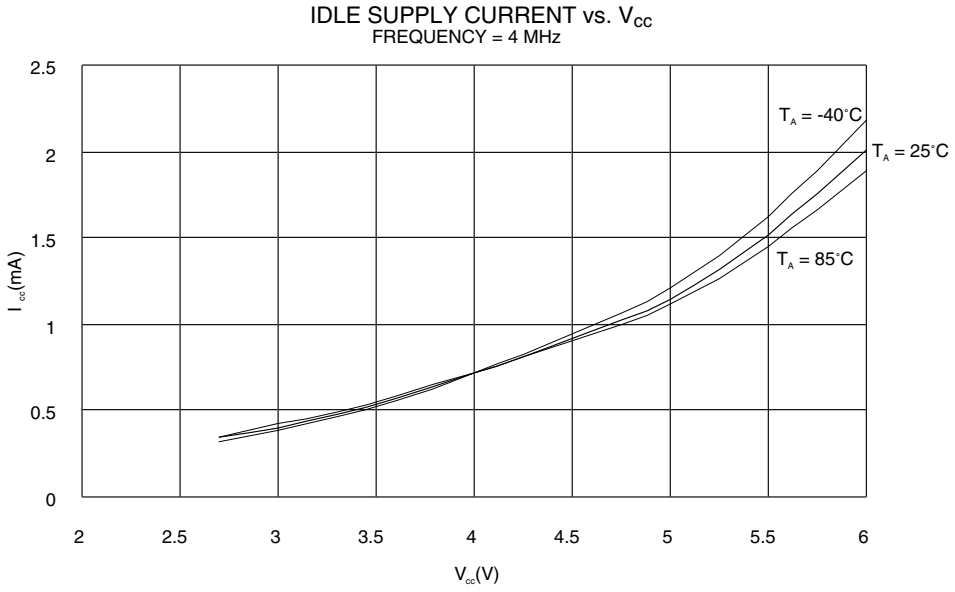


Figure 43. Idle Supply Current vs. V_{CC}, Device Clocked by Internal Oscillator

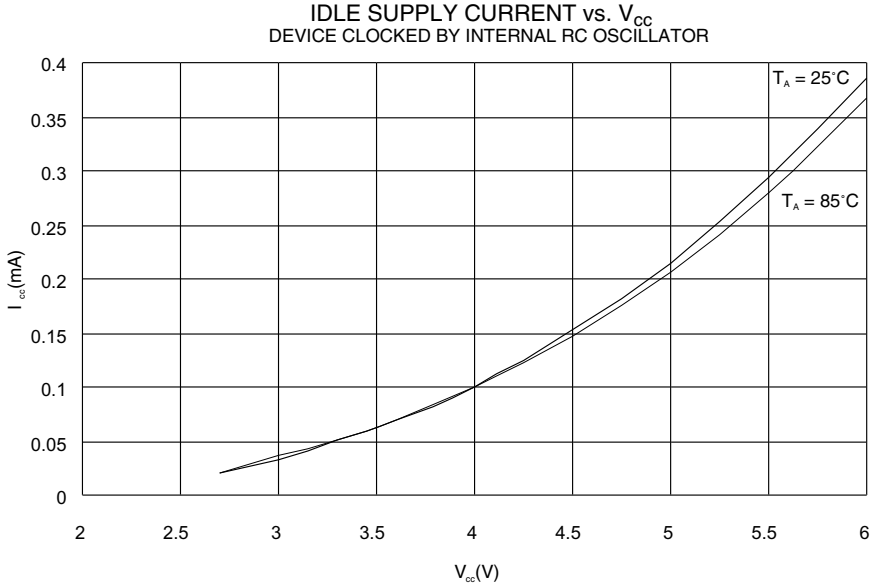


Figure 44. Power-down Supply Current vs. V_{CC}, Watchdog Timer Disabled

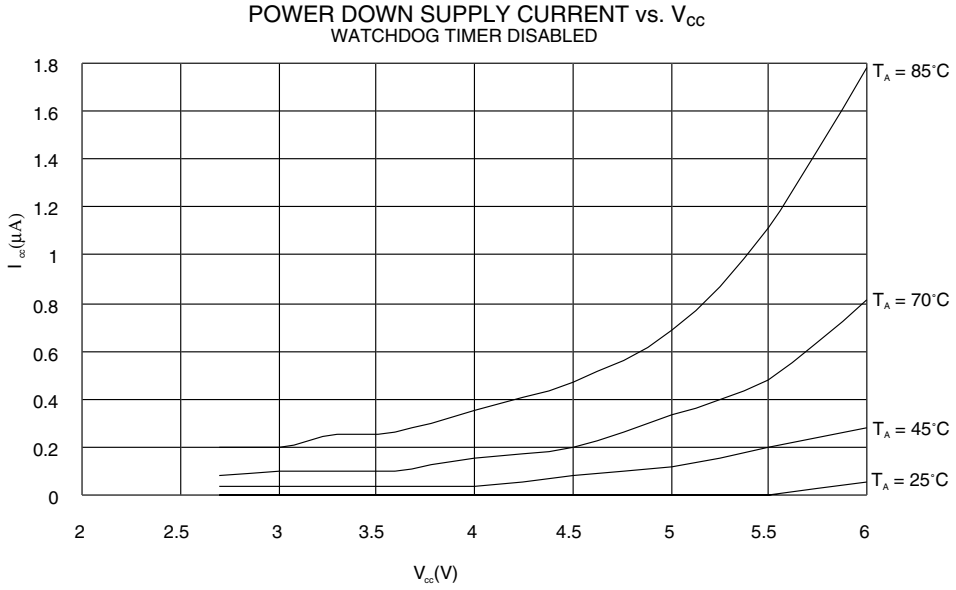


Figure 45. Power-down Supply Current vs. V_{CC}, Watchdog Timer Enabled

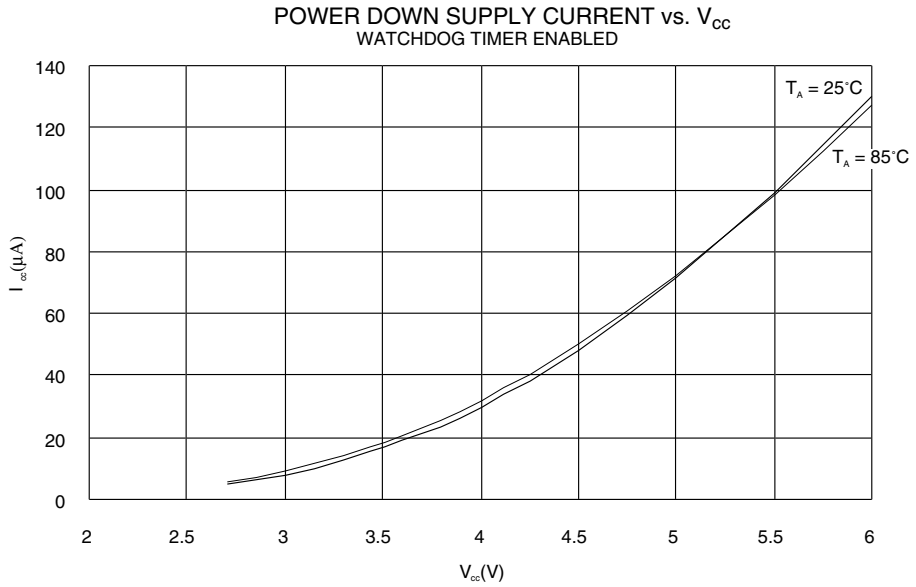


Figure 46. Internal RC Oscillator Frequency vs. V_{CC}

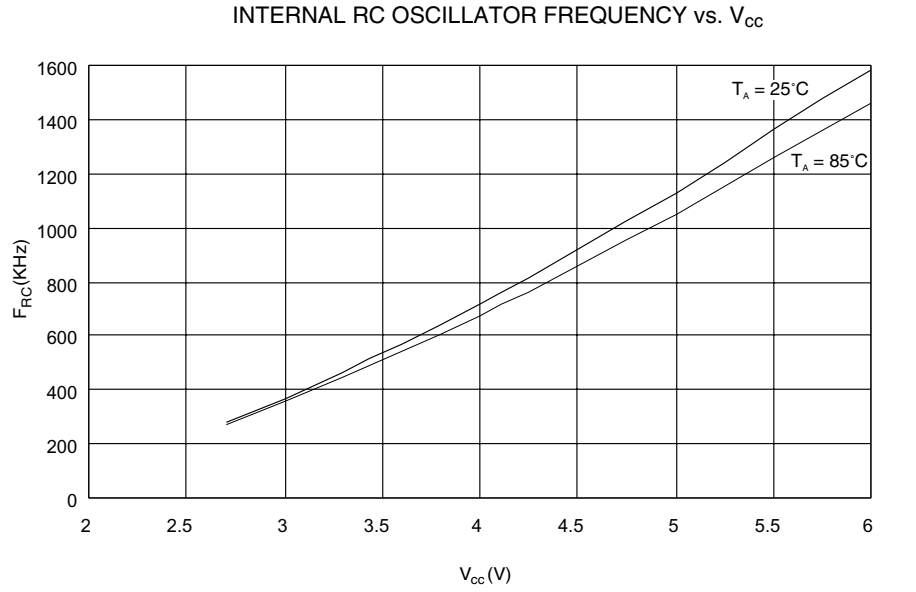
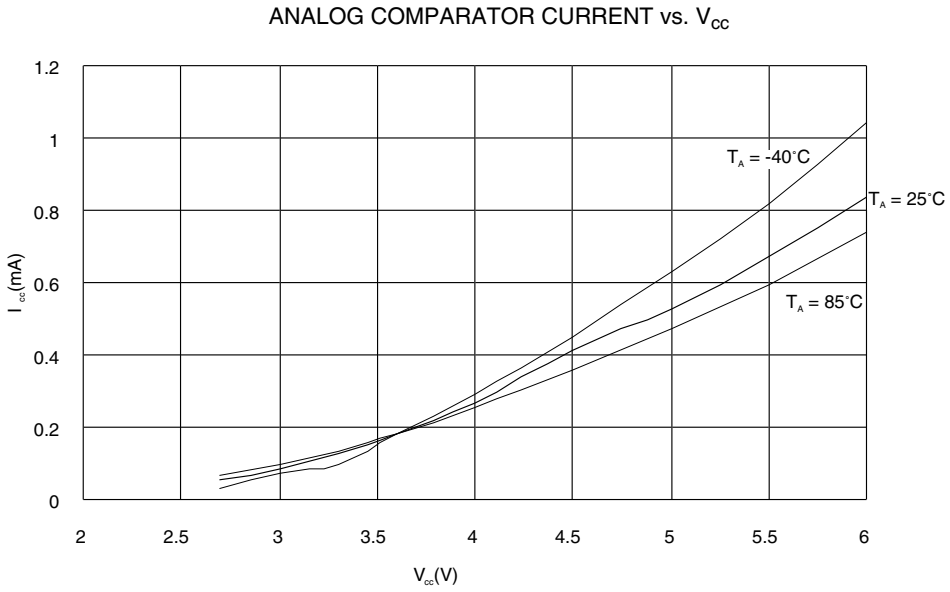


Figure 47. Analog Comparator Current vs. V_{CC}



Note: Analog comparator offset voltage is measured as absolute offset.

Figure 48. Analog Comparator Offset Voltage vs. Common Mode Voltage

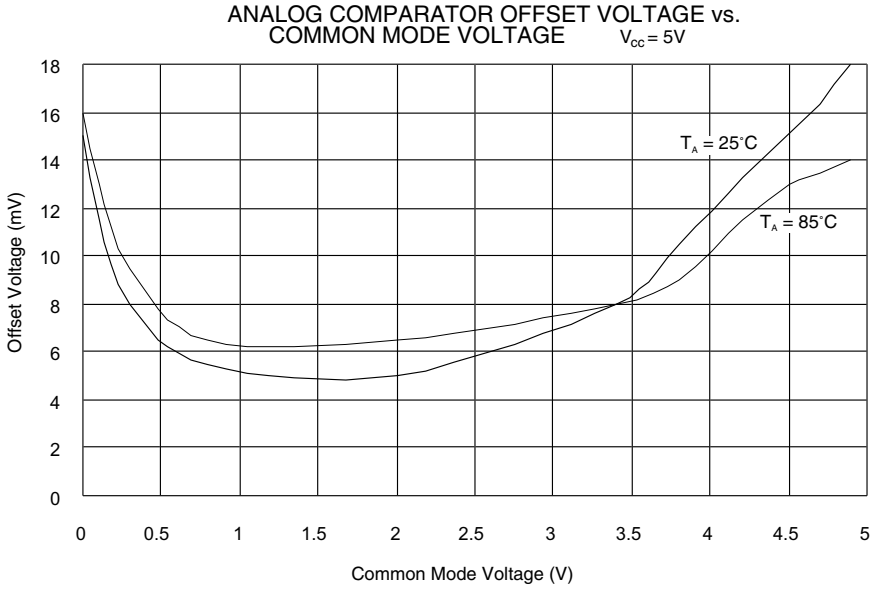


Figure 49. Analog Comparator Offset Voltage vs. Common Mode Voltage

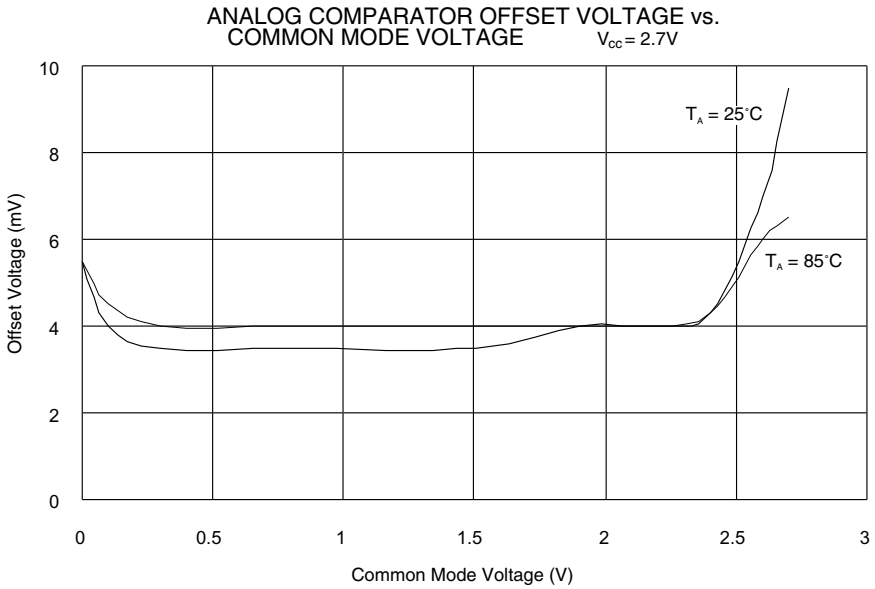
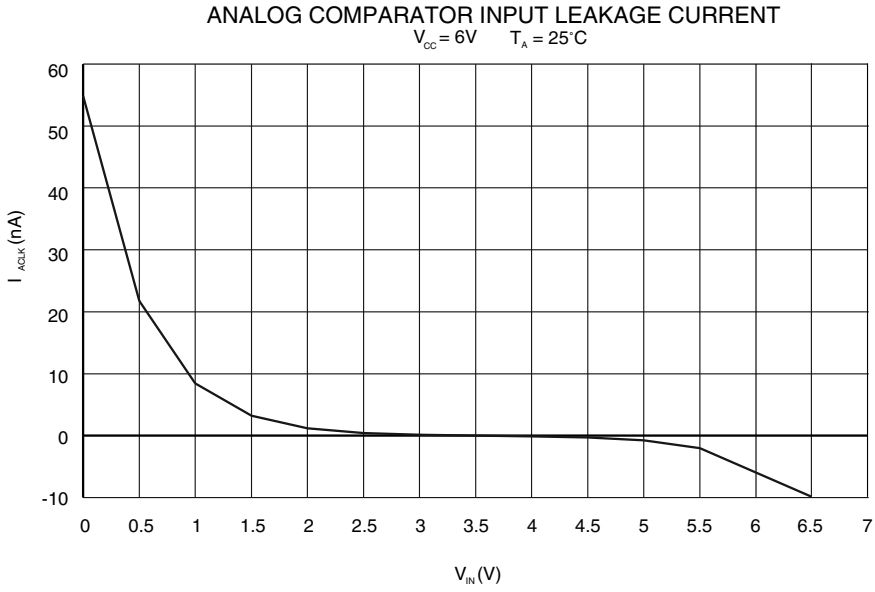


Figure 50. Analog Comparator Input Leakage Current



Note: Sink and source capabilities of I/O ports are measured on one pin at a time.

Figure 51. Pull-up Resistor Current vs. Input Voltage

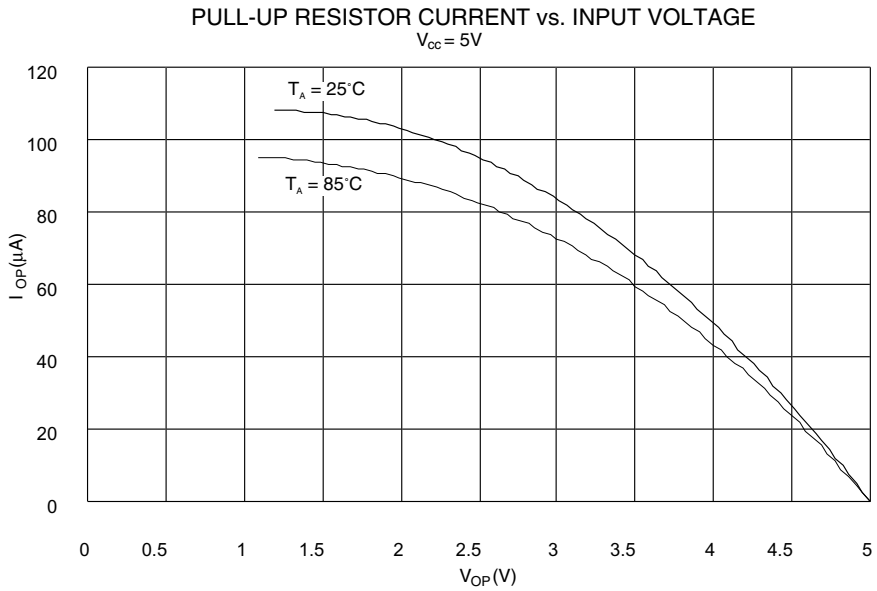


Figure 52. Pull-up Resistor Current vs. Input Voltage

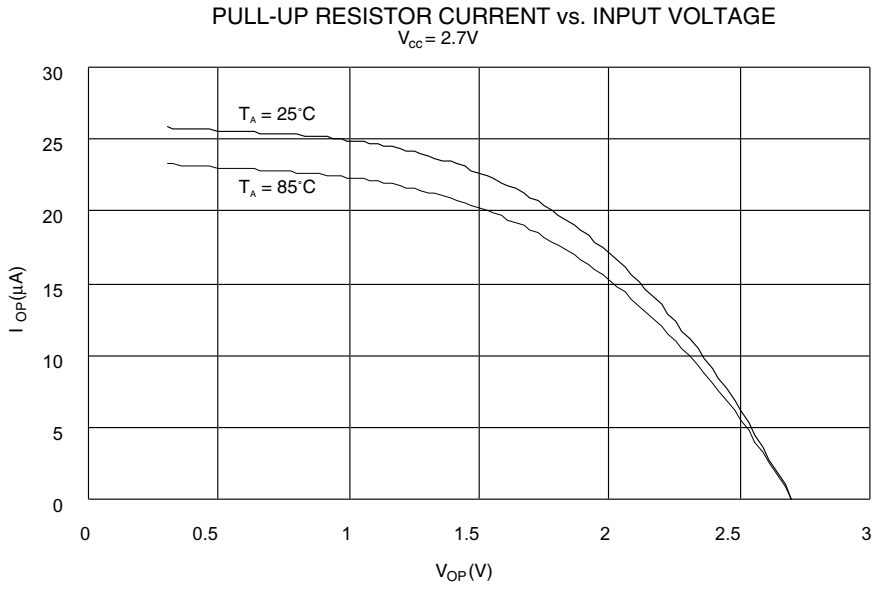


Figure 53. I/O Pin Sink Current vs. Output Voltage

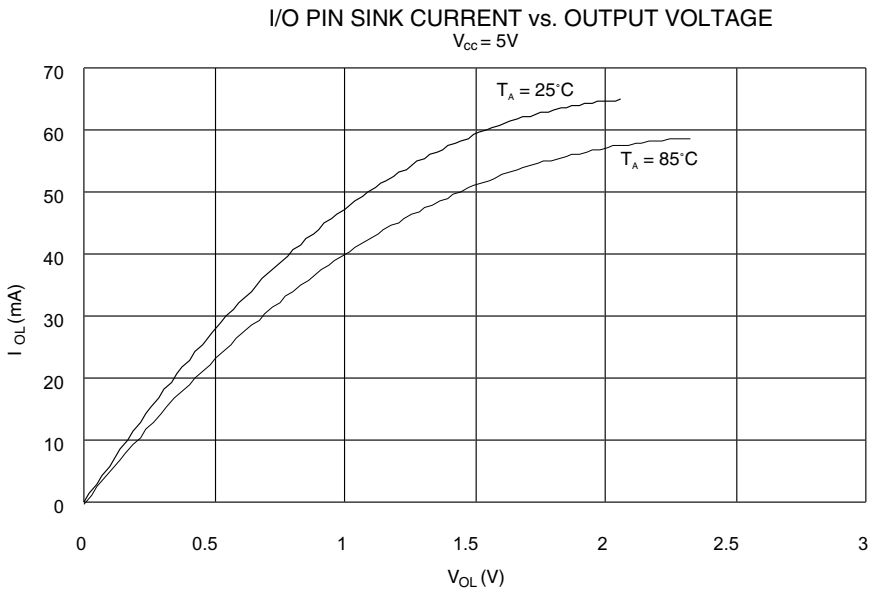


Figure 54. I/O Pin Source Current vs. Output Voltage

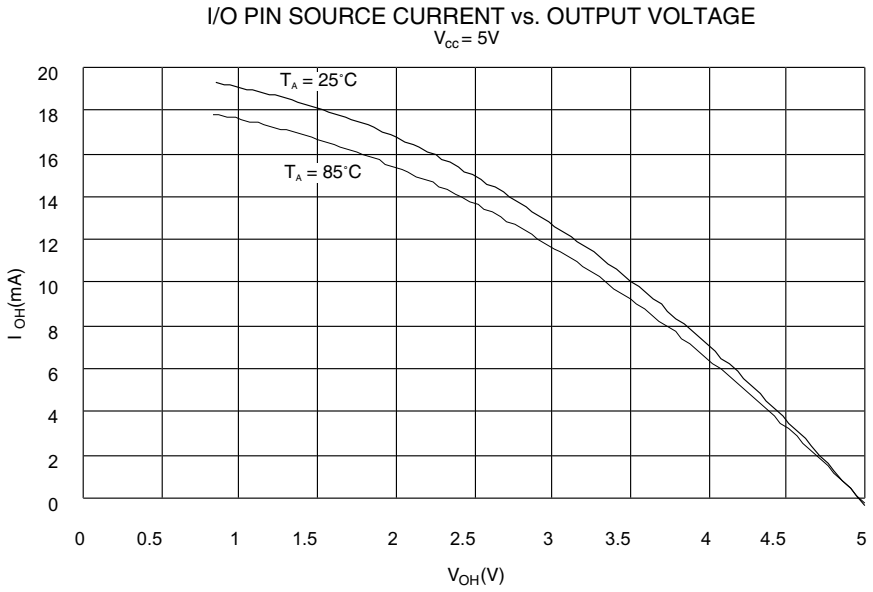


Figure 55. I/O Pin Sink Current vs. Output Voltage

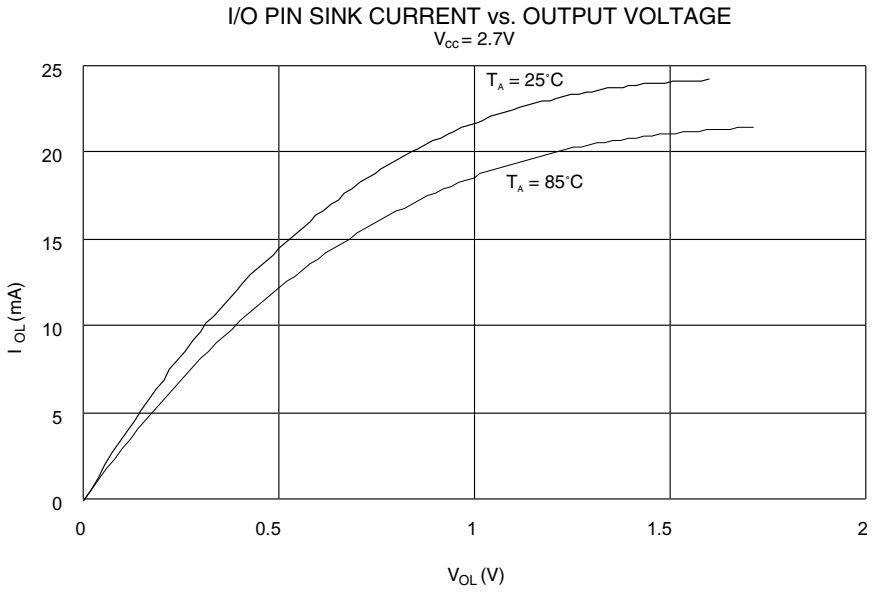
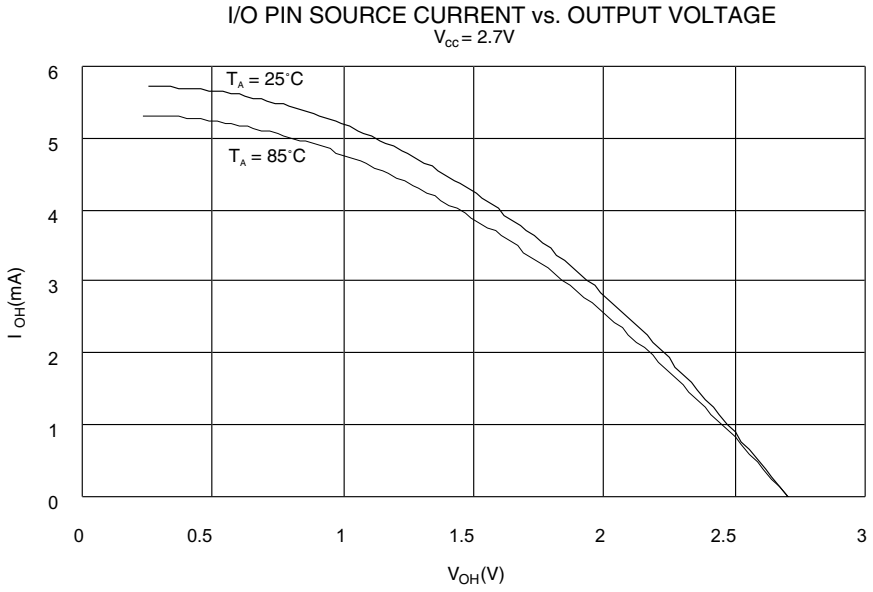


Figure 56. I/O Pin Source Current vs. Output Voltage



Note: Input threshold is measured at the center point of the hysteresis.

Figure 57. I/O Pin Input Threshold Voltage vs. V_{CC}

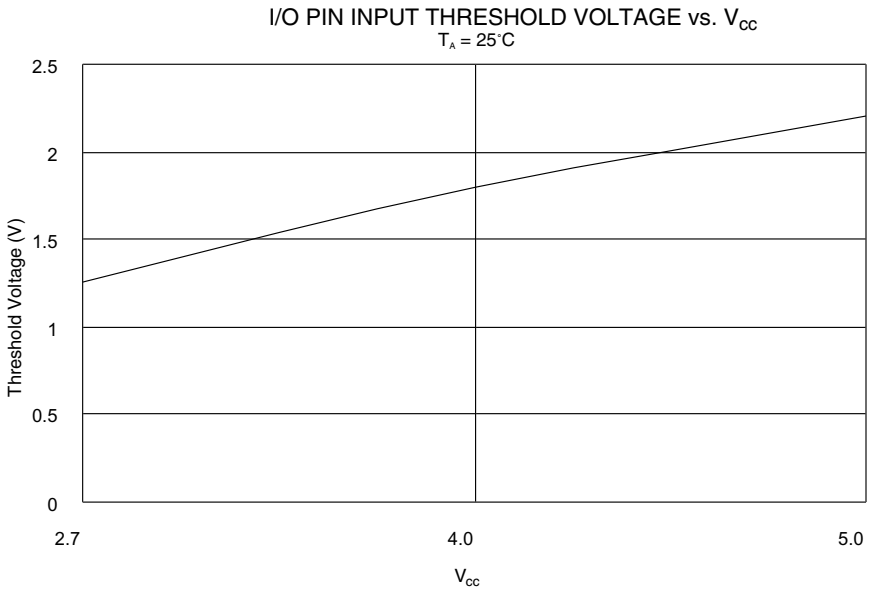
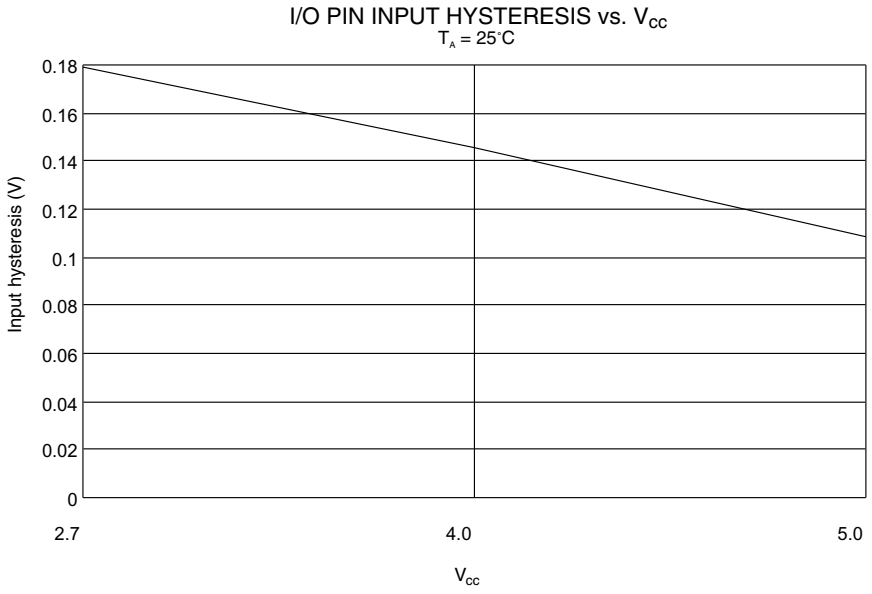


Figure 58. I/O Pin Input Hysteresis vs. V_{CC}



AT90S1200 Register Summary

Address	Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Page
$3F	SREG	I	T	H	S	V	N	Z	C	page 11
$3E	Reserved									
$3D	Reserved									
$3C	Reserved									
$3B	GIMSK	-	INT0	-	-	-	-	-	-	page 15
$3A	Reserved									
$39	TIMSK	-	-	-	-	-	-	TOIE0	-	page 16
$38	TIFR	-	-	-	-	-	-	TOV0	-	page 16
$37	Reserved									
$36	Reserved									
$35	MCUCR	-	-	SE	SM	-	-	ISC01	ISC00	page 18
$34	Reserved									
$33	TCCR0	-	-	-	-	-	CS02	CS01	CS00	page 21
$32	TCNT0	Timer/Counter0 (8 Bits)								page 22
$31	Reserved									
$30	Reserved									
$2F	Reserved									
$2E	Reserved									
$2D	Reserved									
$2C	Reserved									
$2B	Reserved									
$2A	Reserved									
$29	Reserved									
$28	Reserved									
$27	Reserved									
$26	Reserved									
$25	Reserved									
$24	Reserved									
$23	Reserved									
$22	Reserved									
$21	WDTCR	-	-	-	-	WDE	WDP2	WDP1	WDP0	page 23
$20	Reserved									
$1F	Reserved									
$1E	EEAR	-	EEPROM Address Register							page 25
$1D	EEDR	EEPROM Data Register								page 25
$1C	EECR	-	-	-	-	-	-	EEWE	EERE	page 25
$1B	Reserved									
$1A	Reserved									
$19	Reserved									
$18	PORTB	PORTB7	PORTB6	PORTB5	PORTB4	PORTB3	PORTB2	PORTB1	PORTB0	page 29
$17	DDRB	DDB7	DDB6	DDB5	DDB4	DDB3	DDB2	DDB1	DDB0	page 29
$16	PINB	PINB7	PINB6	PINB5	PINB4	PINB3	PINB2	PINB1	PINB0	page 29
$15	Reserved									
$14	Reserved									
$13	Reserved									
$12	PORTD	-	PORTD6	PORTD5	PORTD4	PORTD3	PORTD2	PORTD1	PORTD0	page 34
$11	DDRD	-	DDD6	DDD5	DDD4	DDD3	DDD2	DDD1	DDD0	page 34
$10	PIND	-	PIND6	PIND5	PIND4	PIND3	PIND2	PIND1	PIND0	page 34
$0F	Reserved									
...	Reserved									
$09	Reserved									
$08	ACSR	ACD	-	ACO	ACI	ACIE	-	ACIS1	ACIS0	page 27
…	Reserved									
$00	Reserved									

Notes: 1. For compatibility with future devices, reserved bits should be written to zero if accessed. Reserved I/O memory addresses should never be written.

2. Some of the status flags are cleared by writing a logical “1” to them. Note that the CBI and SBI instructions will operate on all bits in the I/O register, writing a “1” back into any flag read as set, thus clearing the flag. The CBI and SBI instructions work with registers $00 to $1F only.

Instruction Set Summary

Mnemonic	Operands	Description	Operation	Flags	# Clocks
ARITHMETIC AND LOGIC INSTRUCTIONS					
ADD	Rd, Rr	Add Two Registers	Rd ← Rd + Rr	Z,C,N,V,H	1
ADC	Rd, Rr	Add with Carry Two Registers	Rd ← Rd + Rr + C	Z,C,N,V,H	1
SUB	Rd, Rr	Subtract Two Registers	Rd ← Rd - Rr	Z,C,N,V,H	1
SUBI	Rd, K	Subtract Constant from Register	Rd ← Rd - K	Z,C,N,V,H	1
SBC	Rd, Rr	Subtract with Carry Two Registers	Rd ← Rd - Rr - C	Z,C,N,V,H	1
SBCI	Rd, K	Subtract with Carry Constant from Reg.	Rd ← Rd - K - C	Z,C,N,V,H	1
AND	Rd, Rr	Logical AND Registers	Rd ← Rd • Rr	Z,N,V	1
ANDI	Rd, K	Logical AND Register and Constant	Rd ← Rd • K	Z,N,V	1
OR	Rd, Rr	Logical OR Registers	Rd ← Rd v Rr	Z,N,V	1
ORI	Rd, K	Logical OR Register and Constant	Rd ← Rd v K	Z,N,V	1
EOR	Rd, Rr	Exclusive OR Registers	Rd ← Rd ⊕ Rr	Z,N,V	1
COM	Rd	One's Complement	Rd ← $FF - Rd	Z,C,N,V	1
NEG	Rd	Two's Complement	Rd ← $00 - Rd	Z,C,N,V,H	1
SBR	Rd, K	Set Bit(s) in Register	Rd ← Rd v K	Z,N,V	1
CBR	Rd, K	Clear Bit(s) in Register	Rd ← Rd • (FFh - K)	Z,N,V	1
INC	Rd	Increment	Rd ← Rd + 1	Z,N,V	1
DEC	Rd	Decrement	Rd ← Rd - 1	Z,N,V	1
TST	Rd	Test for Zero or Minus	Rd ← Rd • Rd	Z,N,V	1
CLR	Rd	Clear Register	Rd ← Rd ⊕ Rd	Z,N,V	1
SER	Rd	Set Register	Rd ← $FF	None	1
BRANCH INSTRUCTIONS					
RJMP	k	Relative Jump	PC ← PC + k + 1	None	2
RCALL	k	Relative Subroutine Call	PC ← PC + k + 1	None	3
RET		Subroutine Return	PC ← STACK	None	4
RETI		Interrupt Return	PC ← STACK	I	4
CPSE	Rd, Rr	Compare, Skip if Equal	if (Rd = Rr) PC ← PC + 2 or 3	None	1/2
CP	Rd, Rr	Compare	Rd - Rr	Z,N,V,C,H	1
CPC	Rd, Rr	Compare with Carry	Rd - Rr - C	Z,N,V,C,H	1
CPI	Rd, K	Compare Register with Immediate	Rd - K	Z,N,V,C,H	1
SBRC	Rr, b	Skip if Bit in Register Cleared	if (Rr(b) = 0) PC ← PC + 2 or 3	None	1/2
SBRS	Rr, b	Skip if Bit in Register is Set	if (Rr(b) = 1) PC ← PC + 2 or 3	None	1/2
SBIC	P, b	Skip if Bit in I/O Register Cleared	if (P(b)= 0) PC ← PC + 2 or 3	None	1/2
SBIS	P, b	Skip if Bit in I/O Register is Set	if (P(b) = 1) PC ← PC + 2 or 3	None	1/2
BRBS	s, k	Branch if Status Flag Set	if (SREG(s) = 1) then PC ← PC + k + 1	None	1/2
BRBC	s, k	Branch if Status Flag Cleared	if (SREG(s) = 0) then PC ← PC + k + 1	None	1/2
BREQ	k	Branch if Equal	if (Z = 1) then PC ← PC + k + 1	None	1/2
BRNE	k	Branch if Not Equal	if (Z = 0) then PC ← PC + k + 1	None	1/2
BRCS	k	Branch if Carry Set	if (C = 1) then PC ← PC + k + 1	None	1/2
BRCC	k	Branch if Carry Cleared	if (C = 0) then PC ← PC + k + 1	None	1/2
BRSH	k	Branch if Same or Higher	if (C = 0) then PC ← PC + k + 1	None	1/2
BRLO	k	Branch if Lower	if (C = 1) then PC ← PC + k + 1	None	1/2
BRMI	k	Branch if Minus	if (N = 1) then PC ← PC + k + 1	None	1/2
BRPL	k	Branch if Plus	if (N = 0) then PC ← PC + k + 1	None	1/2
BRGE	k	Branch if Greater or Equal, Signed	if (N ⊕ V = 0) then PC ← PC + k + 1	None	1/2
BRLT	k	Branch if Less than Zero, Signed	if (N ⊕ V = 1) then PC ← PC + k + 1	None	1/2
BRHS	k	Branch if Half-carry Flag Set	if (H = 1) then PC ← PC + k + 1	None	1/2
BRHC	k	Branch if Half-carry Flag Cleared	if (H = 0) then PC ← PC + k + 1	None	1/2
BRTS	k	Branch if T-Flag Set	if (T = 1) then PC ← PC + k + 1	None	1/2
BRTC	k	Branch if T-Flag Cleared	if (T = 0) then PC ← PC + k + 1	None	1/2
BRVS	k	Branch if Overflow Flag is Set	if (V = 1) then PC ← PC + k + 1	None	1/2
BRVC	k	Branch if Overflow Flag is Cleared	if (V = 0) then PC ← PC + k + 1	None	1/2
BRIE	k	Branch if Interrupt Enabled	if (I = 1) then PC ← PC + k + 1	None	1/2
BRID	k	Branch if Interrupt Disabled	if (I = 0) then PC ← PC + k + 1	None	1/2
DATA TRANSFER INSTRUCTIONS					
LD	Rd, Z	Load Register Indirect	Rd ← (Z)	None	2
ST	Z, Rr	Store Register Indirect	(Z) ← Rr	None	2
MOV	Rd, Rr	Move between Registers	Rd ← Rr	None	1
LDI	Rd, K	Load Immediate	Rd ← K	None	1
IN	Rd, P	In Port	Rd ← P	None	1
OUT	P, Rr	Out Port	P ← Rr	None	1

Instruction Set Summary (Continued)

Mnemonic	Operands	Description	Operation	Flags	# Clocks
BIT AND BIT-TEST INSTRUCTIONS					
SBI	P, b	Set Bit in I/O Register	I/O(P,b) ← 1	None	2
CBI	P, b	Clear Bit in I/O Register	I/O(P,b) ← 0	None	2
LSL	Rd	Logical Shift Left	Rd(n+1) ← Rd(n), Rd(0) ← 0	Z,C,N,V	1
LSR	Rd	Logical Shift Right	Rd(n) ← Rd(n+1), Rd(7) ← 0	Z,C,N,V	1
ROL	Rd	Rotate Left through Carry	Rd(0) ← C,Rd(n+1) ← Rd(n),C ← Rd(7)	Z,C,N,V	1
ROR	Rd	Rotate Right through Carry	Rd(7) ← C,Rd(n) ← Rd(n+1),C ← Rd(0)	Z,C,N,V	1
ASR	Rd	Arithmetic Shift Right	Rd(n) ← Rd(n+1), n = 0..6	Z,C,N,V	1
SWAP	Rd	Swap Nibbles	Rd(3..0) ← Rd(7..4),Rd(7..4) ← Rd(3..0)	None	1
BSET	s	Flag Set	SREG(s) ← 1	SREG(s)	1
BCLR	s	Flag Clear	SREG(s) ← 0	SREG(s)	1
BST	Rr, b	Bit Store from Register to T	T ← Rr(b)	T	1
BLD	Rd, b	Bit Load from T to Register	Rd(b) ← T	None	1
SEC		Set Carry	C ← 1	C	1
CLC		Clear Carry	C ← 0	C	1
SEN		Set Negative Flag	N ← 1	N	1
CLN		Clear Negative Flag	N ← 0	N	1
SEZ		Set Zero Flag	Z ← 1	Z	1
CLZ		Clear Zero Flag	Z ← 0	Z	1
SEI		Global Interrupt Enable	I ← 1	I	1
CLI		Global Interrupt Disable	I ← 0	I	1
SES		Set Signed Test Flag	S ← 1	S	1
CLS		Clear Signed Test Flag	S ← 0	S	1
SEV		Set Two's Complement Overflow	V ← 1	V	1
CLV		Clear Two's Complement Overflow	V ← 0	V	1
SET		Set T in SREG	T ← 1	T	1
CLT		Clear T in SREG	T ← 0	T	1
SEH		Set Half-carry Flag in SREG	H ← 1	H	1
CLH		Clear Half-carry Flag in SREG	H ← 0	H	1
NOP		No Operation		None	1
SLEEP		Sleep	(see specific descr. for Sleep function)	None	1
WDR		Watchdog Reset	(see specific descr. for WDR/timer)	None	1

Ordering Information[1]

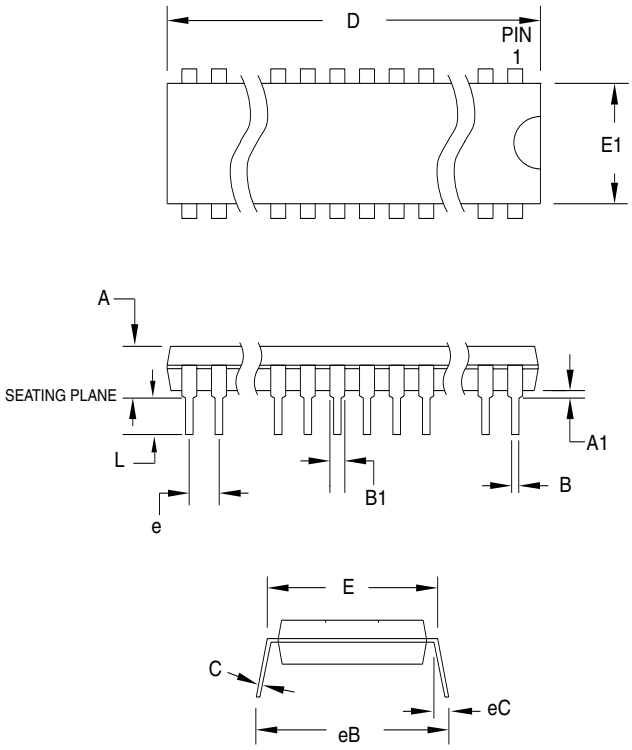
Speed (MHz)	Power Supply	Ordering Code	Package	Operation Range
4	2.7 - 6.0V	AT90S1200-4PC AT90S1200-4SC AT90S1200-4YC	20P3 20S 20Y	Commercial (0°C to 70°C)
		AT90S1200-4PI AT90S1200-4SI AT90S1200-4YI	20P3 20S 20Y	Industrial (-40°C to 85°C)
12	4.0 - 6.0V	AT90S1200-12PC AT90S1200-12SC AT90S1200-12YC	20P3 20S 20Y	Commercial (0°C to 70°C)
		AT90S1200-12PI AT90S1200-12SI AT90S1200-12YI	20P3 20S 20Y	Industrial (-40°C to 85°C)

Note: 1. Order AT90S1200A-XXX for devices with the RCEN Fuse programmed.

Package Type	
20P3	20-lead, 0.300" Wide, Plastic Dual Inline Package (PDIP)
20S	20-lead, 0.300" Wide, Plastic Gull Wing Small Outline (SOIC)
20Y	20-lead, 5.3 mm Wide, Plastic Shrink Small Outline Package (SSOP)

Packaging Information

20P3



COMMON DIMENSIONS
(Unit of Measure = mm)

SYMBOL	MIN	NOM	MAX	NOTE
A	–	–	5.334	
A1	0.381	–	–	
D	25.984	–	25.493	Note 2
E	7.620	–	8.255	
E1	6.096	–	7.112	Note 2
B	0.356	–	0.559	
B1	1.270	–	1.551	
L	2.921	–	3.810	
C	0.203	–	0.356	
eB	–	–	10.922	
eC	0.000	–	1.524	
e	2.540 TYP			

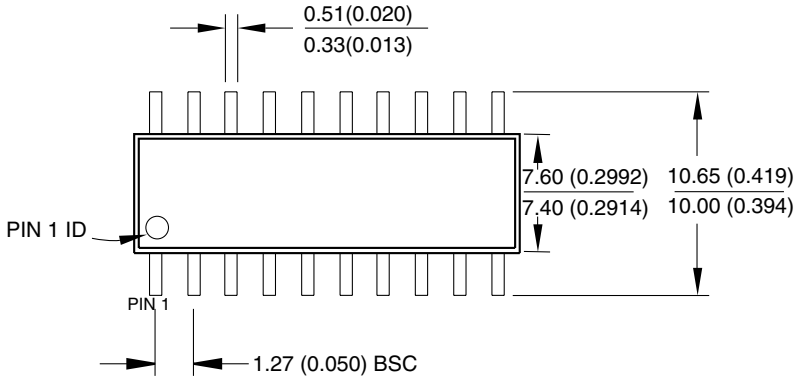
Notes: 1. This package conforms to JEDEC reference MS-001, Variation AD.
2. Dimensions D and E1 do not include mold Flash or Protrusion. Mold Flash or Protrusion shall not exceed 0.25 mm (0.010").

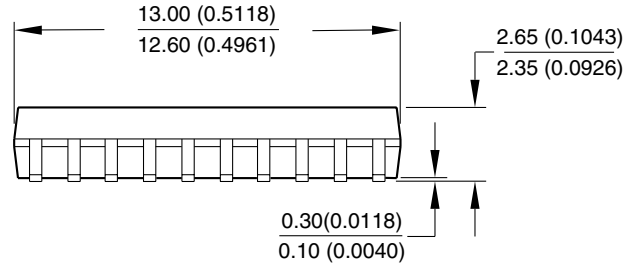
09/28/01

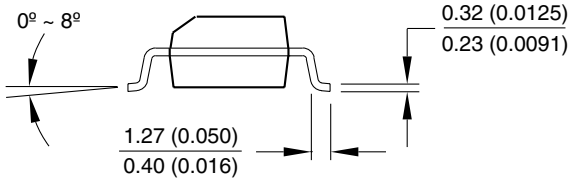
	TITLE	DRAWING NO.	REV.
2325 Orchard Parkway San Jose, CA 95131	**20P3**, 20-lead (0.300"/7.62 mm Wide) Plastic Dual Inline Package (PDIP)	20P3	B

20S

20S, 20-lead, Plastic Gull Wing Small Outline (SOIC), 0.300" body.
Dimensions in Millineters and (Inches)*
JEDEC STANDARD MS-013





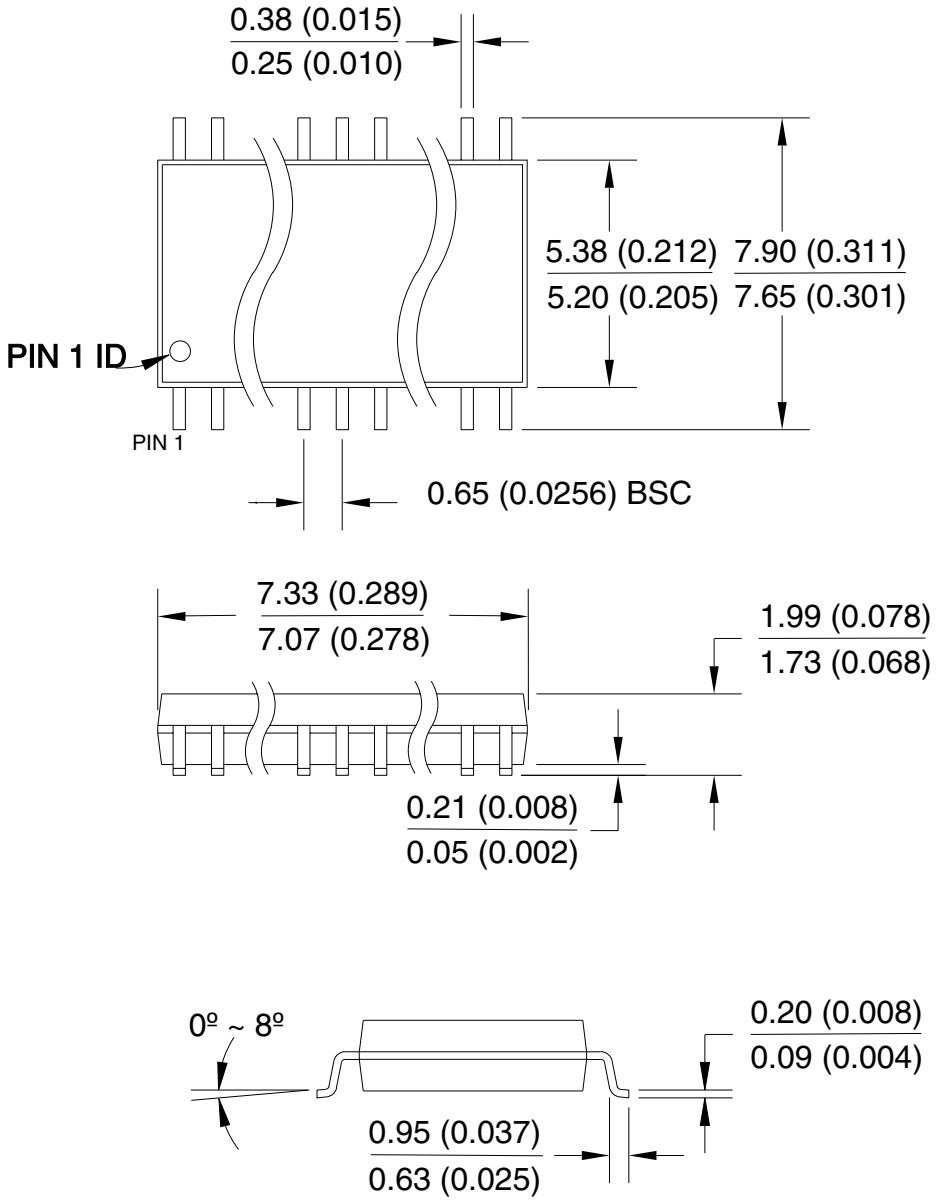


*Controlling dimension: Inches

REV. A 04/11/2001

20Y

20Y, 20-lead Plastic Shrink Small Outline (SSOP), 5.3mm body Width. Dimensions in Millimeters and (inches)*



*Controlling dimension: millimeters

REV. A 04/11/2001

Table of Contents

Features 1

Pin Configuration 1

Description 2

- Block Diagram 2
- Pin Descriptions 3
- Crystal Oscillator 3
- On-chip RC Oscillator 4

Architectural Overview 5

- General Purpose Register File 6
- ALU – Arithmetic Logic Unit 6
- In-System Programmable Flash Program Memory 6
- Program and Data Addressing Modes 7
- Subroutine and Interrupt Hardware Stack 8
- EEPROM Data Memory 9
- Instruction Execution Timing 9
- I/O Memory 10
- Reset and Interrupt Handling 12
- Sleep Modes 19

Timer/Counter0 20

- Timer/Counter0 Prescaler 20

Watchdog Timer 23

EEPROM Read/Write Access 25

- Prevent EEPROM Corruption 26

Analog Comparator 27

I/O Ports 29

- Port B 29
- Port D 34

Memory Programming 37

- Program and Data Memory Lock Bits 37
- Fuse Bits 37
- Signature Bytes 37
- Programming the Flash and EEPROM 37
- Parallel Programming 38
- Parallel Programming Characteristics 43
- Serial Downloading 44
- Serial Programming Characteristics 47

Electrical Characteristics *48*
Absolute Maximum Ratings* 48
DC Characteristics 48
External Clock Drive Waveforms 50
External Clock Drive 50

Typical Characteristics *51*

AT90S1200 Register Summary *62*

Instruction Set Summary *63*

Ordering Information(1) *65*

Packaging Information *66*
20P3 66
20S 67
20Y 68

Table of Contents ..... *i*

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